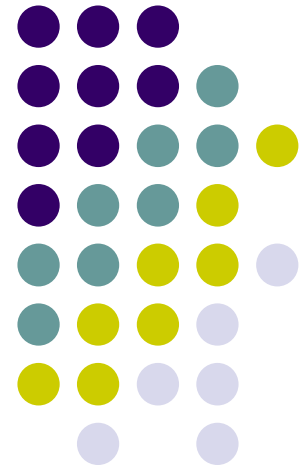


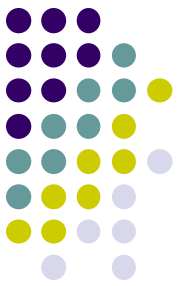


Queues

Data Structures

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Punjabi University

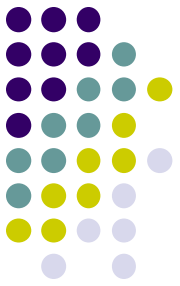


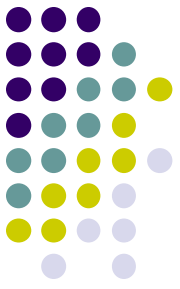


DEFINITION OF QUEUE

- A **Queue** is an ordered collection of items from which items may be deleted at one end (called the **front** of the queue) and into which items may be inserted at the other end (the **rear** of the queue).
- The first element inserted into the queue is the first element to be removed. For this reason a queue is sometimes called a **FIFO** (first-in first-out) list as opposed to the stack, which is a **LIFO** (last-in first-out).

Queues



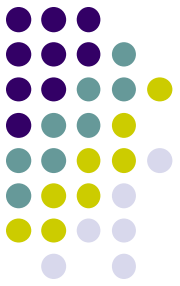


Basic Queue Operations

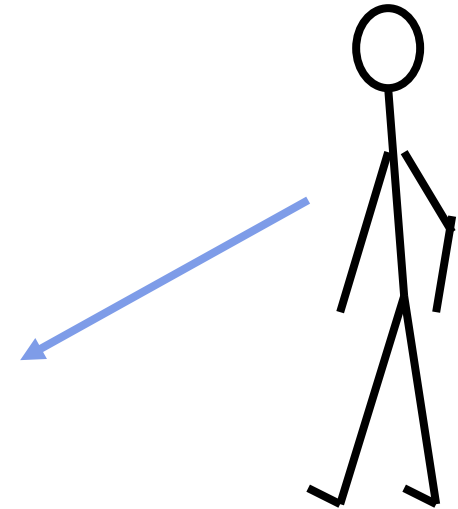
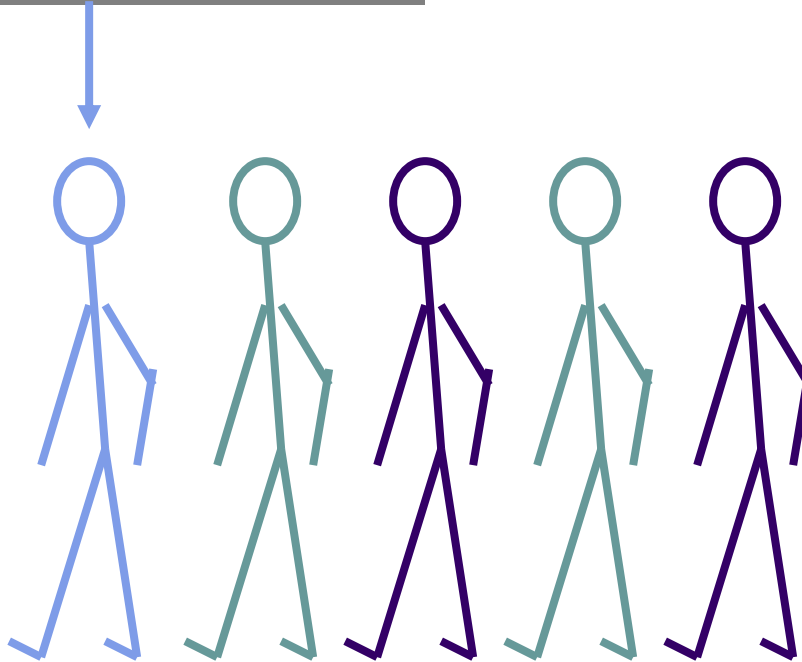
- There are two basic operations that can be performed on a queue:
- **Enqueue:** This operation inserts or adds data in the queue. The data will be added to the rear of the queue.
- **Dequeue:** This operation removes the data, which has been in the queue for the longest time. The data is present in the front of the queue.

Enqueue

Adding an element



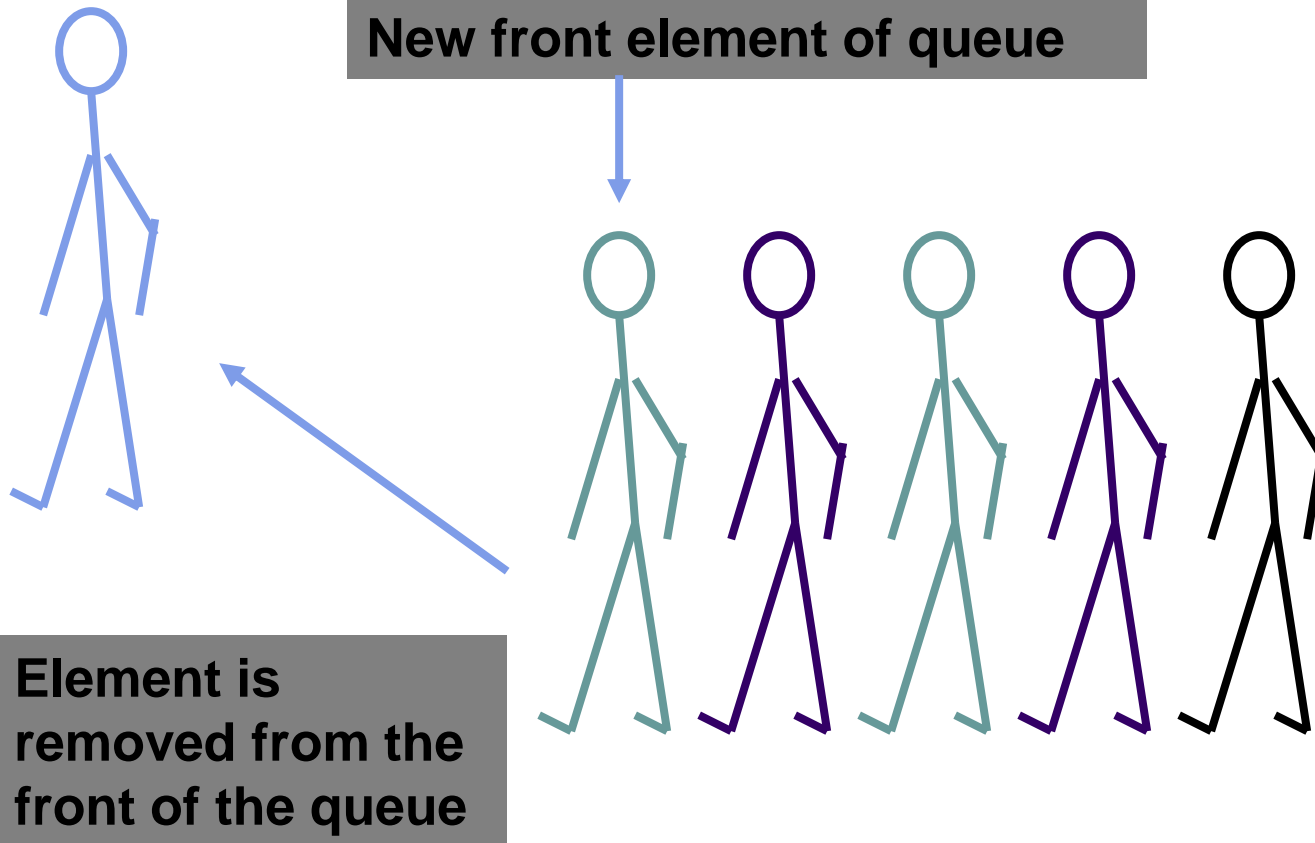
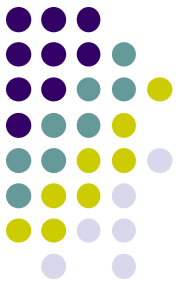
Front of queue

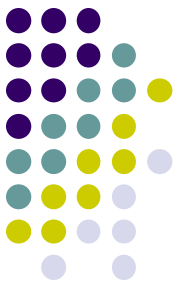


New element is added to the rear of the queue

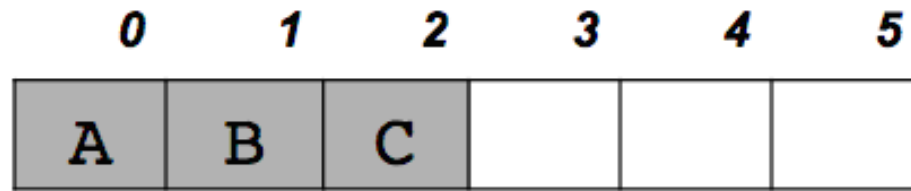
Dequeue

Removing an element





Basic Queue Operations

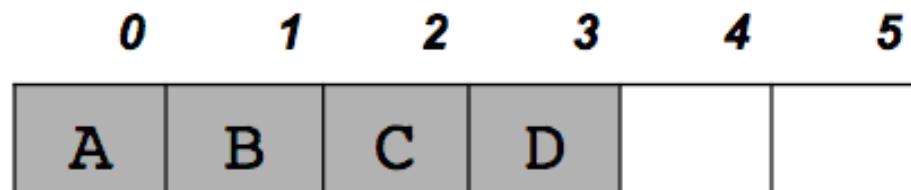


front

rear



enqueue(D)

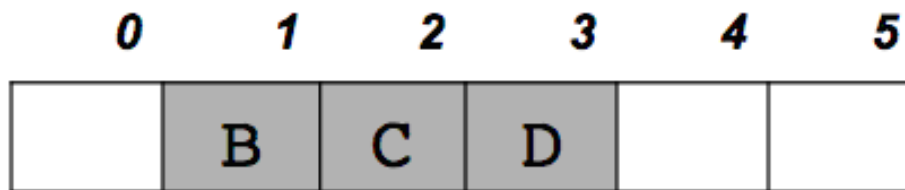


front

rear



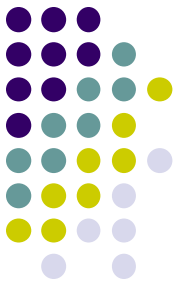
dequeue()



front

rear

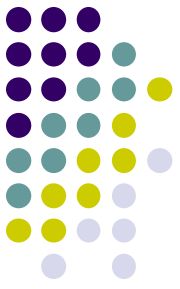
Implementation of the Queues



There are many possible ways to implement the queues:

- An array with front at index 0
- An array with floating front and rear
- A circular array with floating front and rear (wrap around array)
- A singly linked list with front
- A singly linked list with front and rear
- A doubly linked list

Array Implementation of the Queues



- We must keep track of both the *front* and the *rear* of the queue. One method is to keep the *front* of the array in the first location on the array. Then, we can simply increase the counter of the array to show *the rear*.
- Nevertheless, to **delete an entry from this queue is very expensive**, since after the first entry was served, all the existing entry need to be move back one position to fill in the vacancy.
- With a long queue this process can lead to poor performance.

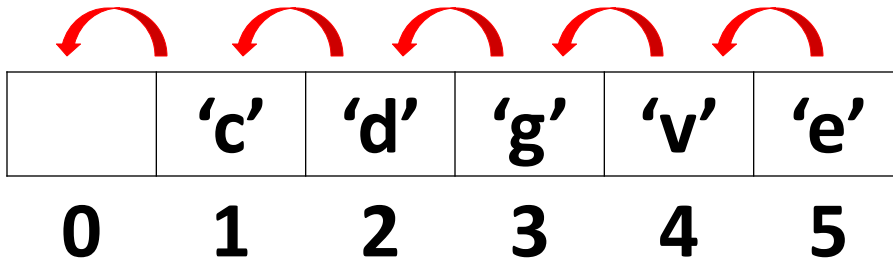


Array with front at index 0

Before:

'a'	'c'	'd'	'g'	'v'	'e'
-----	-----	-----	-----	-----	-----

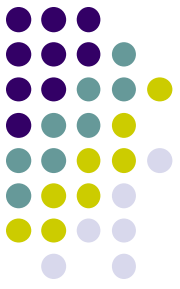
front → 0 1 2 3 4 5 ← *rear*



'a' at index 0 is deleted

After:

'c'	'd'	'g'	'v'	'e'	
0	1	2	3	4	5



Linear Implementation of Queues

- *Indicate the front and rear of the queue.* We can keep track the entry of the queue without moving any entries.
 - Append an entry: increase the rear by one.
 - To remove the entry: increase the front by one.
- *Problem:*
 - Position of front will increase and never decrease
 - This leads to the end of the storage capacity

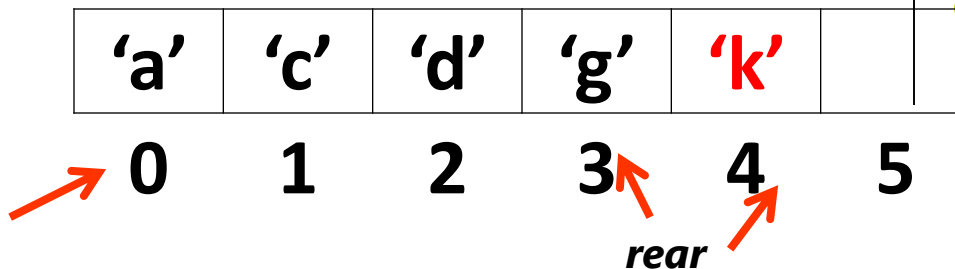
Linear Implementation of Queues



Add 'k' to the queue:

- rear + 1 (increase)

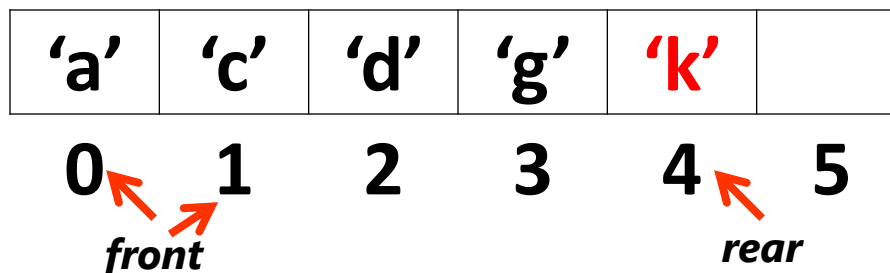
front



Delete 'a' in queue:

- front + 1 (increase)

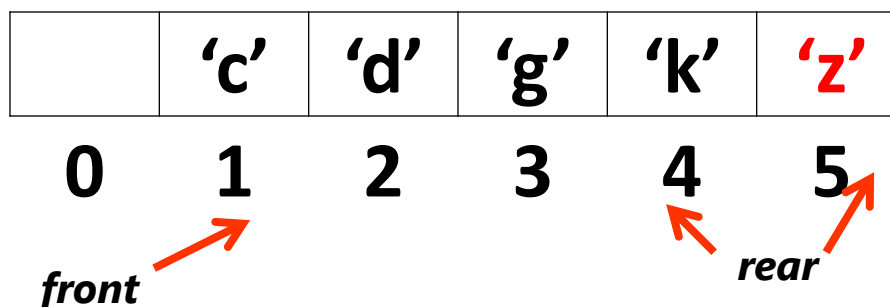
front



Add 'z' to the queue:

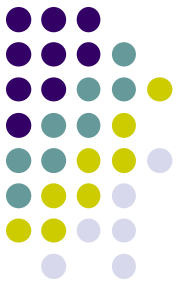
- rear + 1 (increase)

front



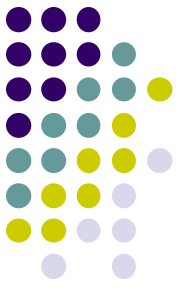
Reach end of storage!!

Need for Circular Queues



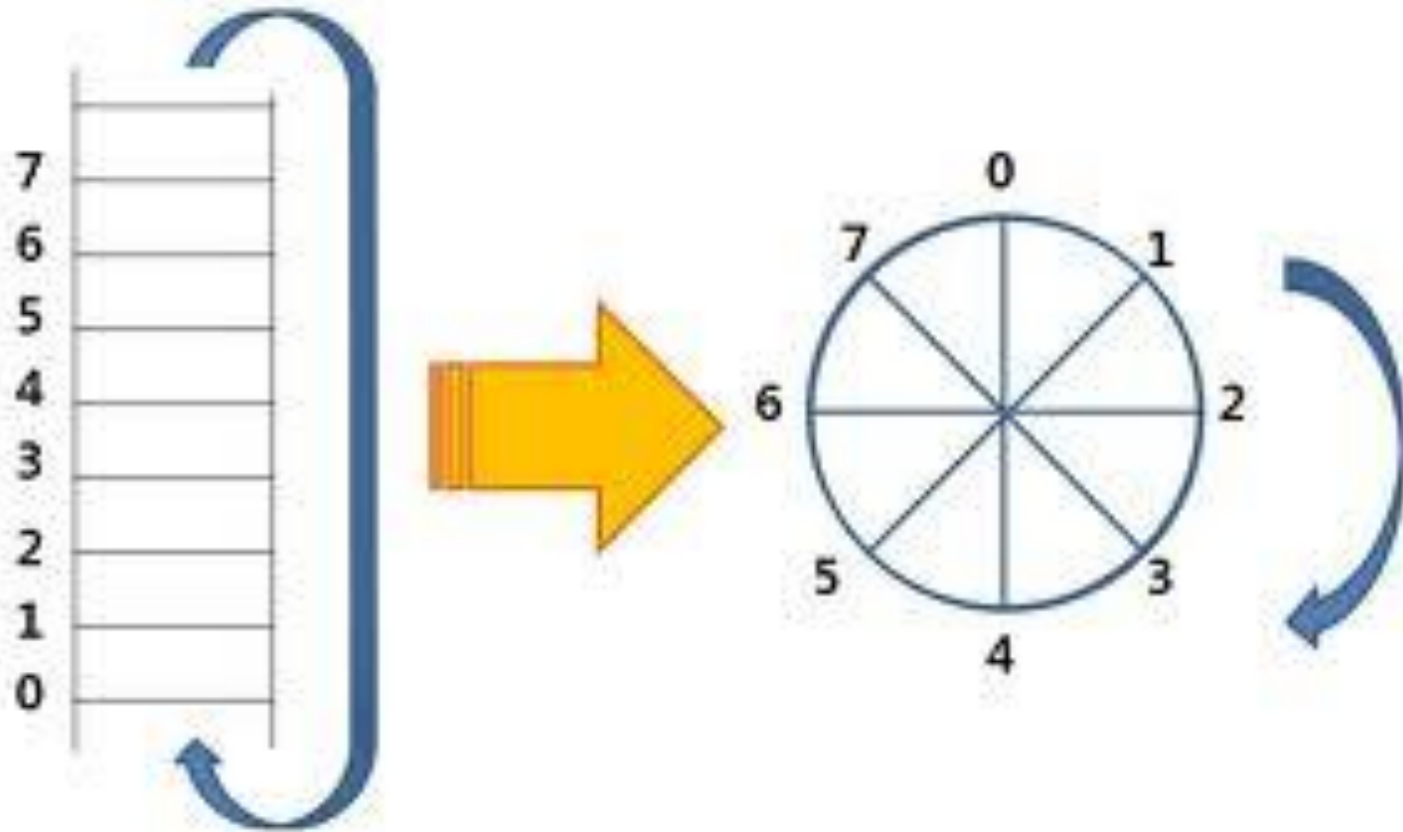
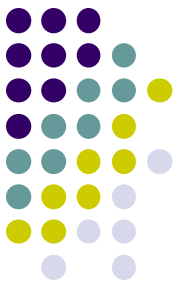
- Queues implemented as linear arrays have the drawback that once the queue is FULL, even though we delete few elements from the "front" and relieve some occupied space, we are not able to add anymore elements, as the "rear" has already reached the Queue's rear most position.

Need for Circular Queues

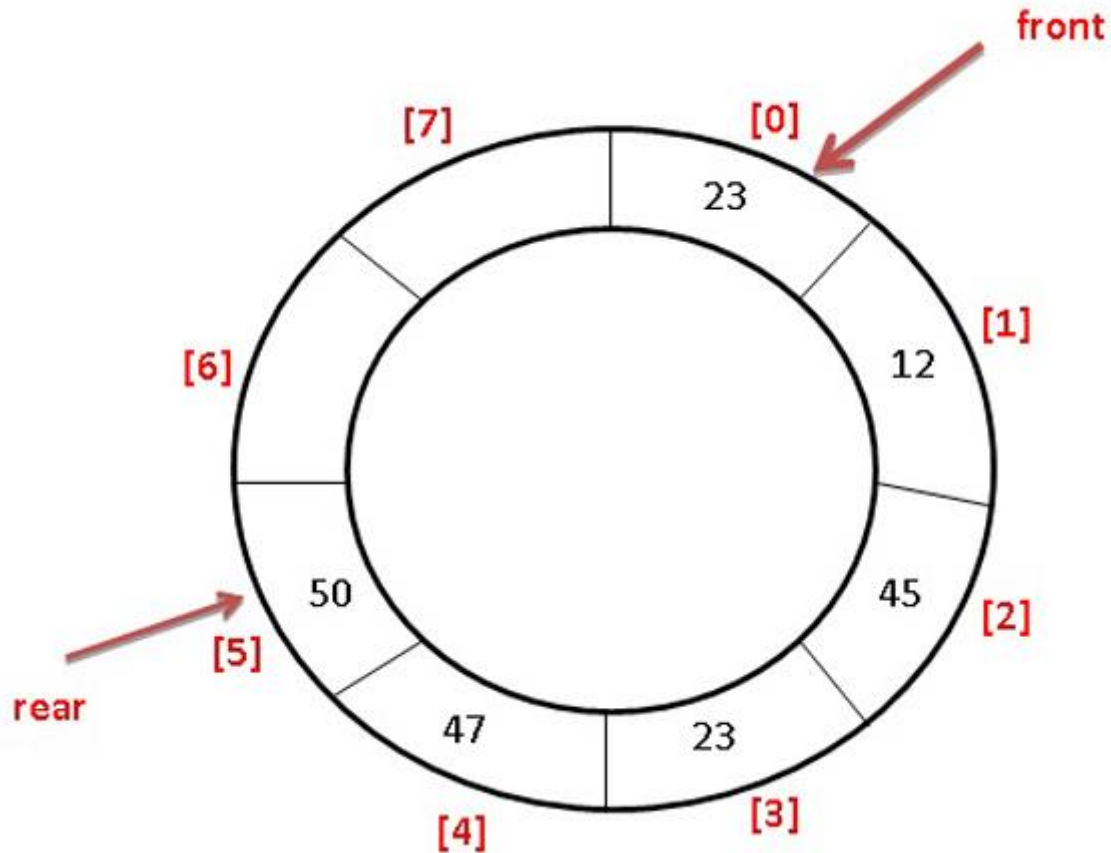
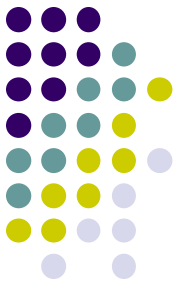


- To solve this problem, queues implement wrapping around. Such queues are called *Circular Queues*.
- Both the front and the rear pointers wrap around to the beginning of the array.
- Circular queues are better than normal queues as they effectively utilise the memory space.

Circular Queues

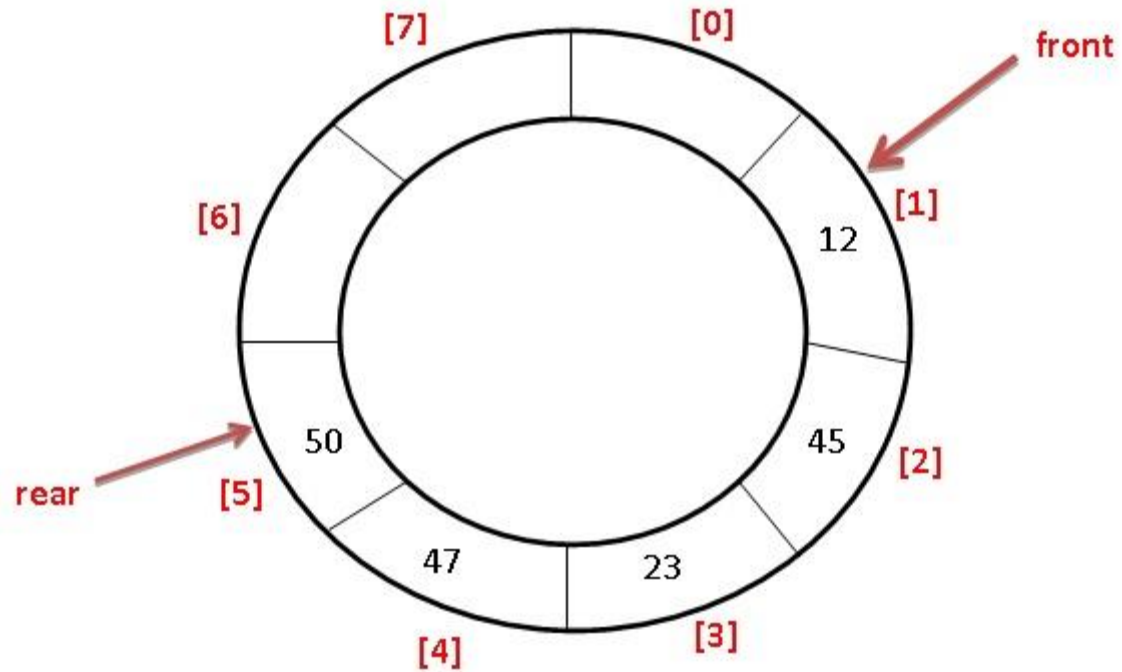
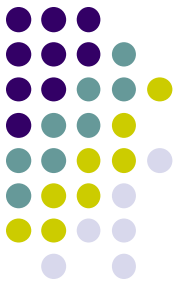


Circular Queues



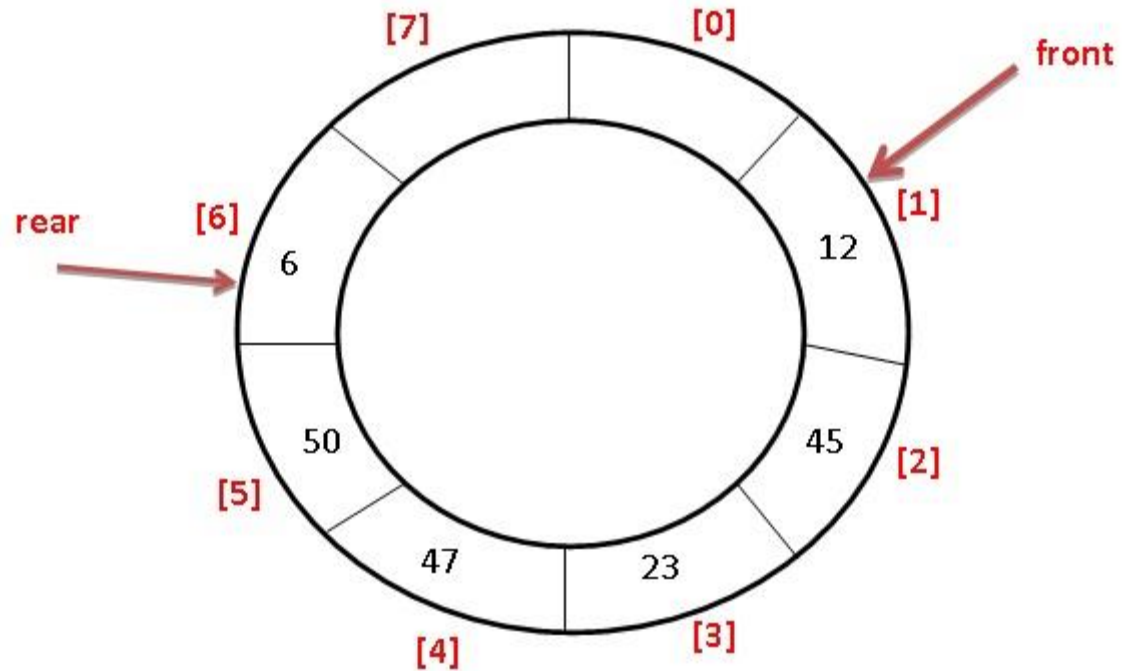
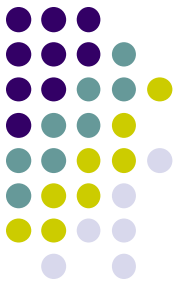
Queue with 6 elements

Circular Queues



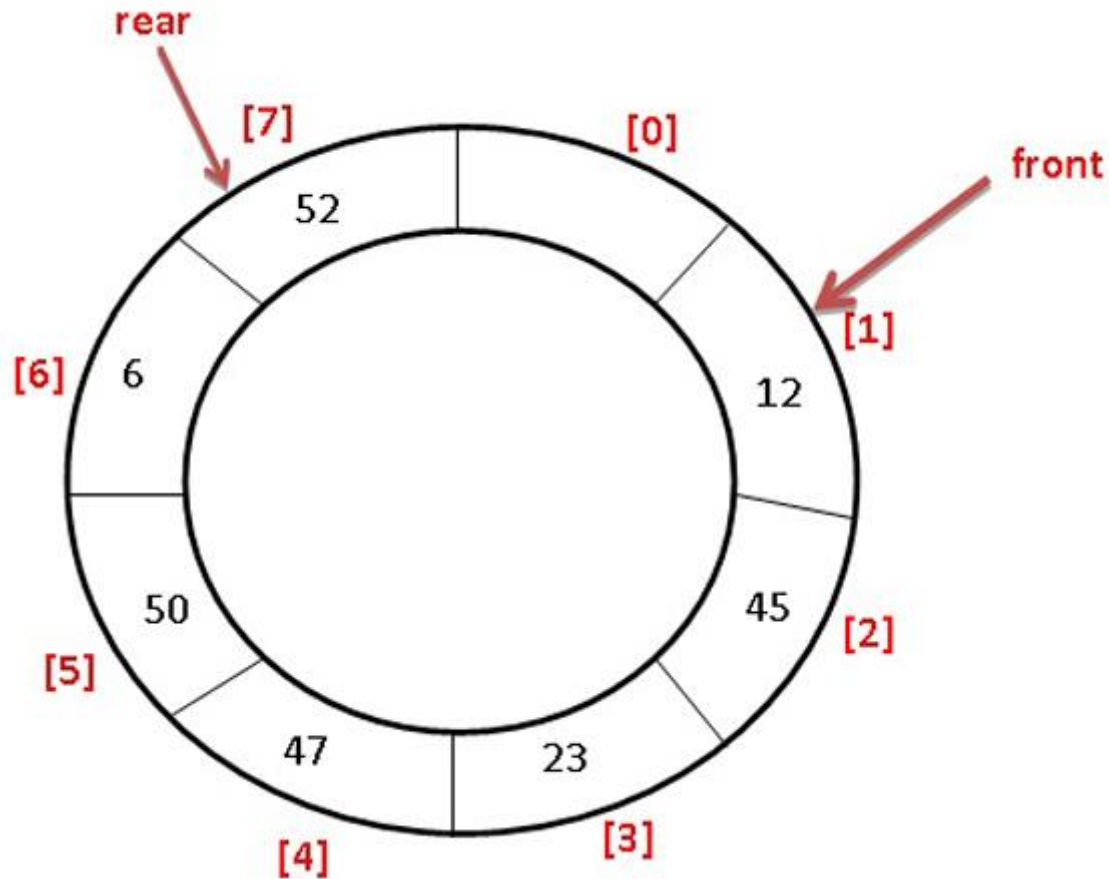
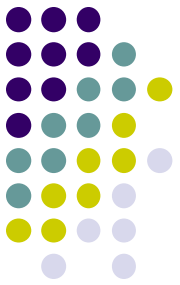
dequeue()

Circular Queues



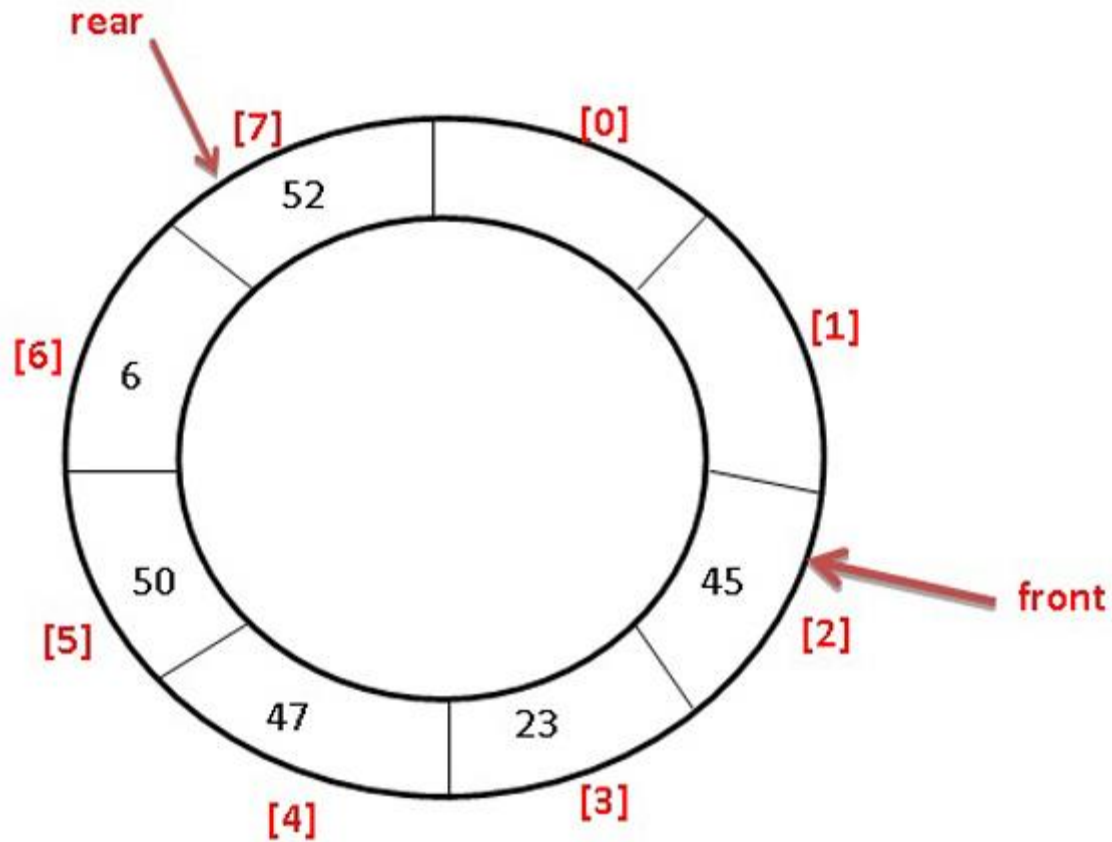
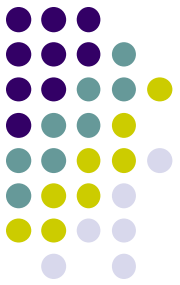
enqueue(6)

Circular Queues



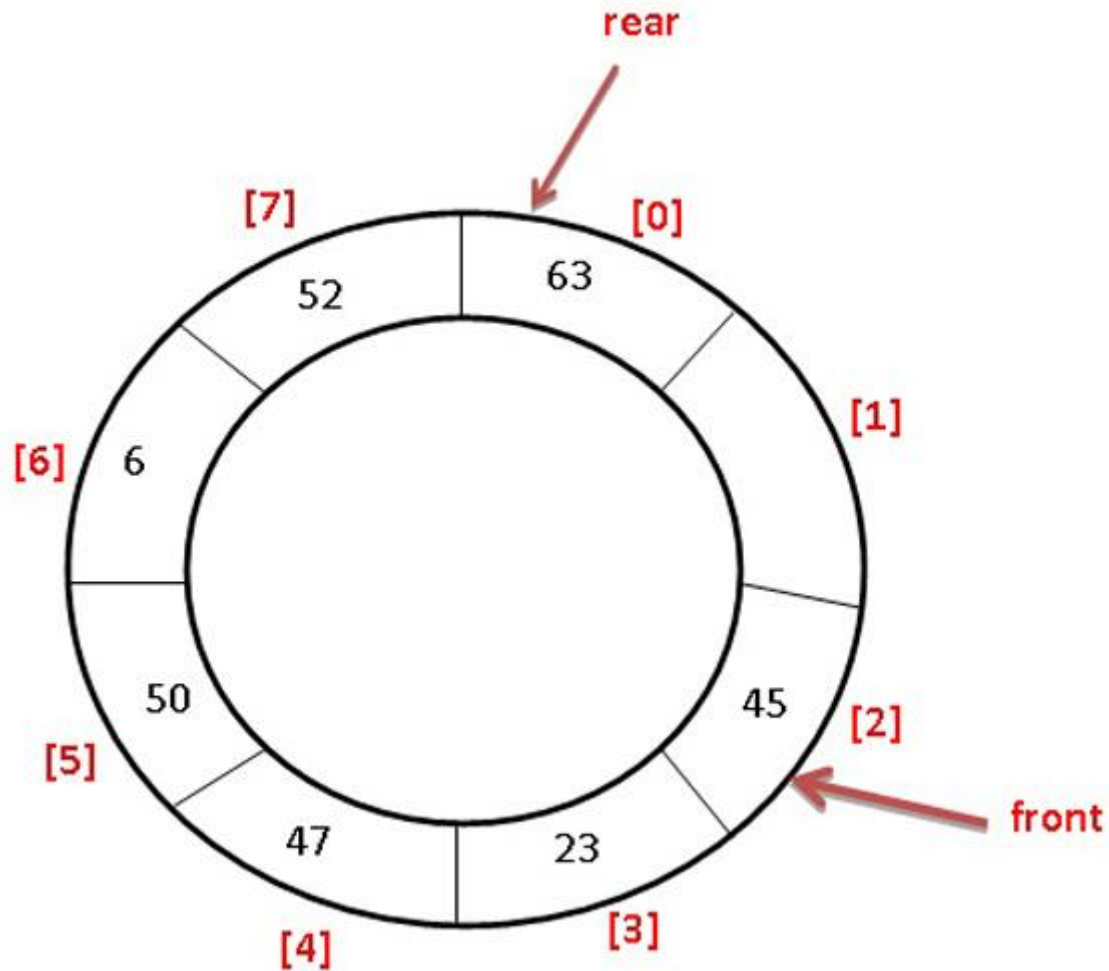
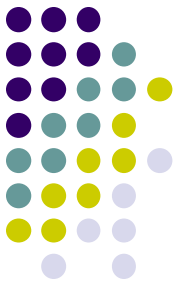
enqueue(52)

Circular Queues



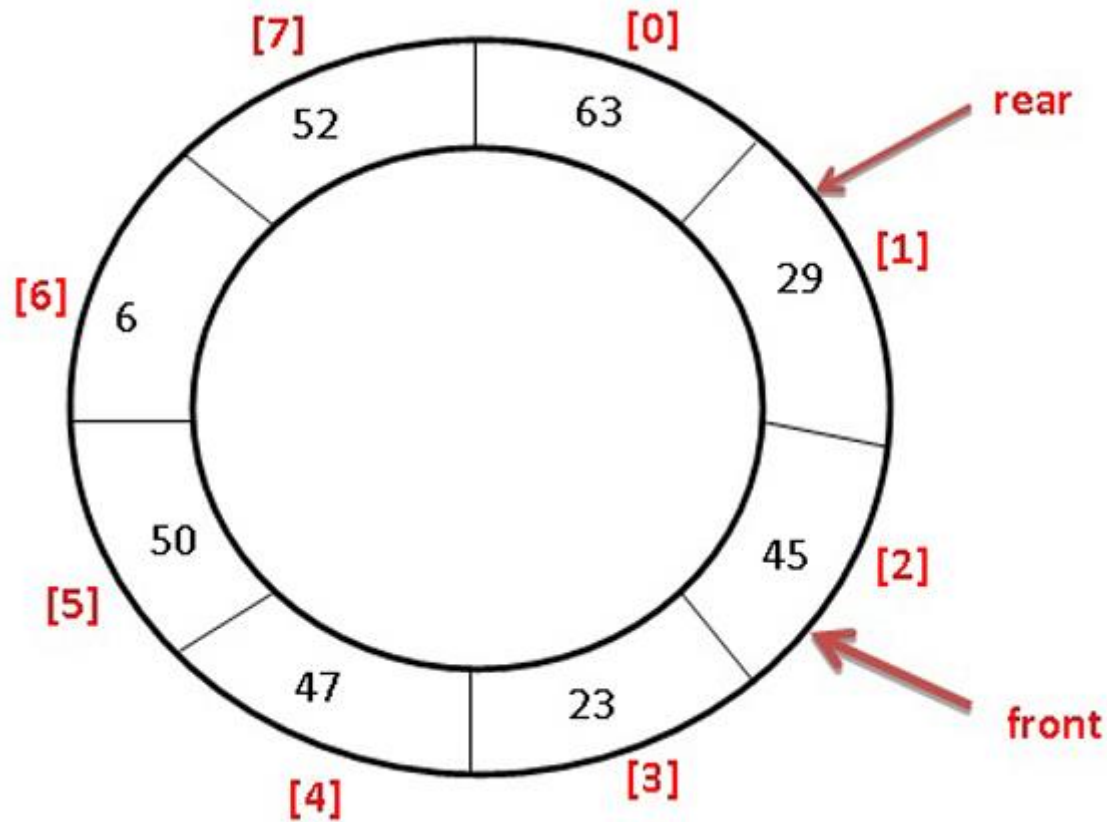
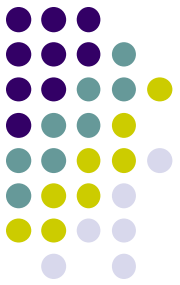
dequeue()

Circular Queues



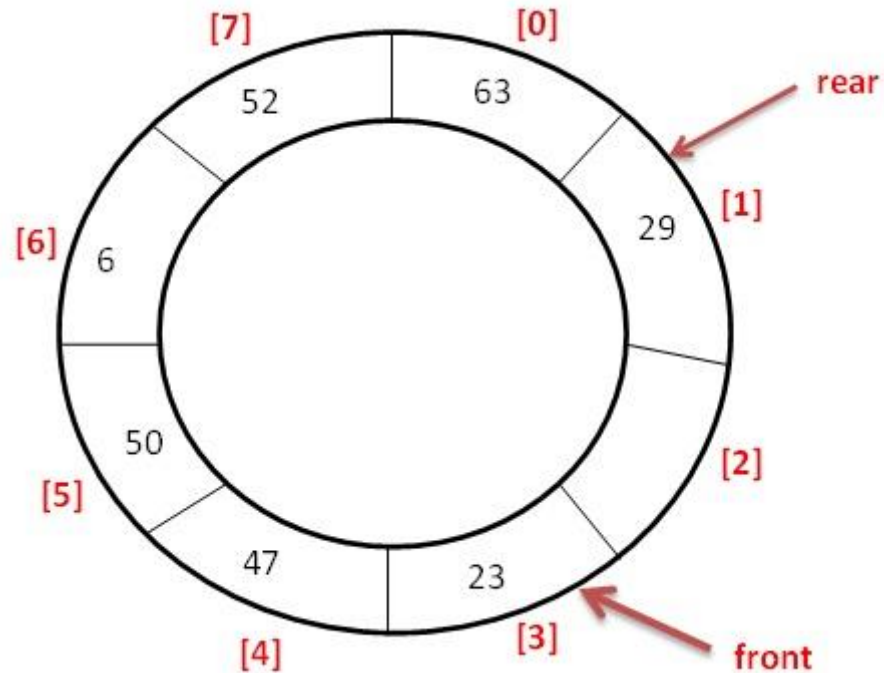
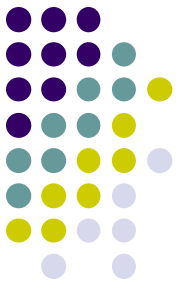
enqueue(63)

Circular Queues



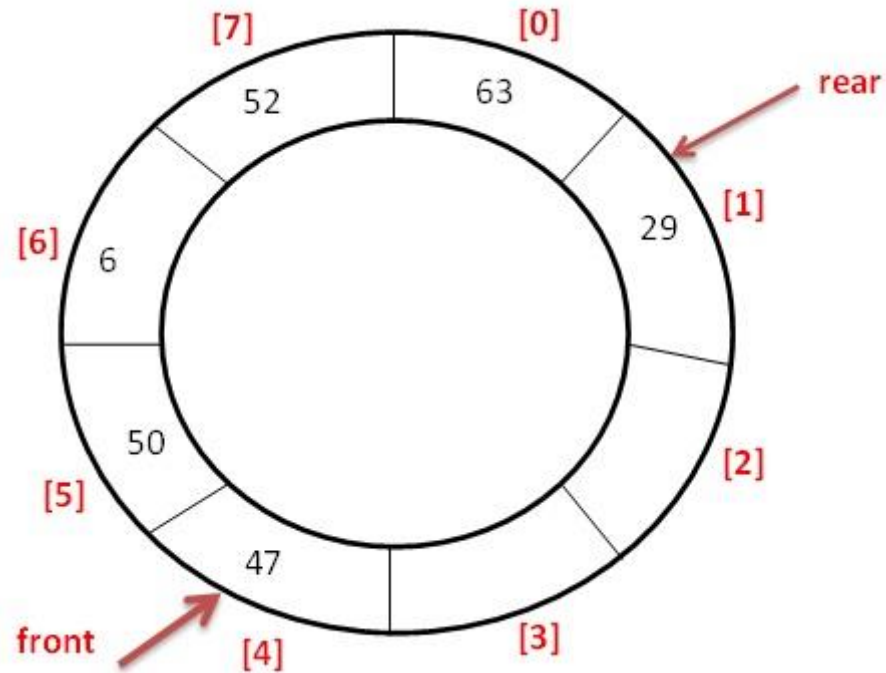
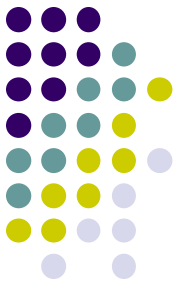
enqueue(29)

Circular Queues



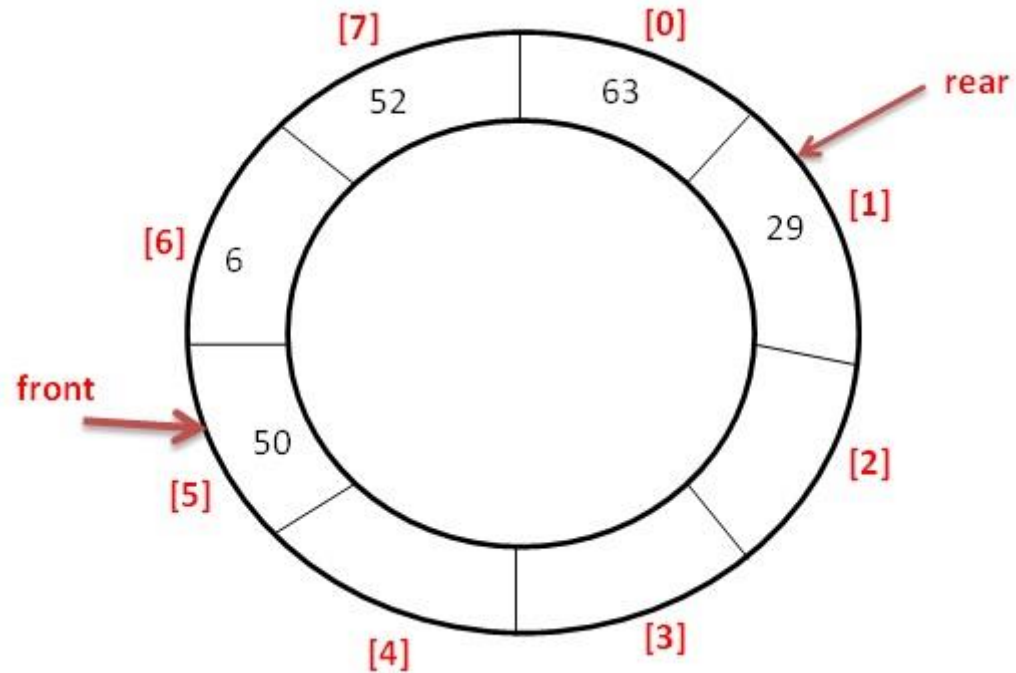
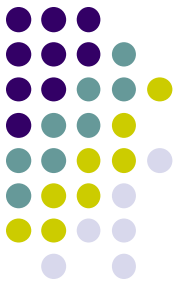
dequeue()

Circular Queues



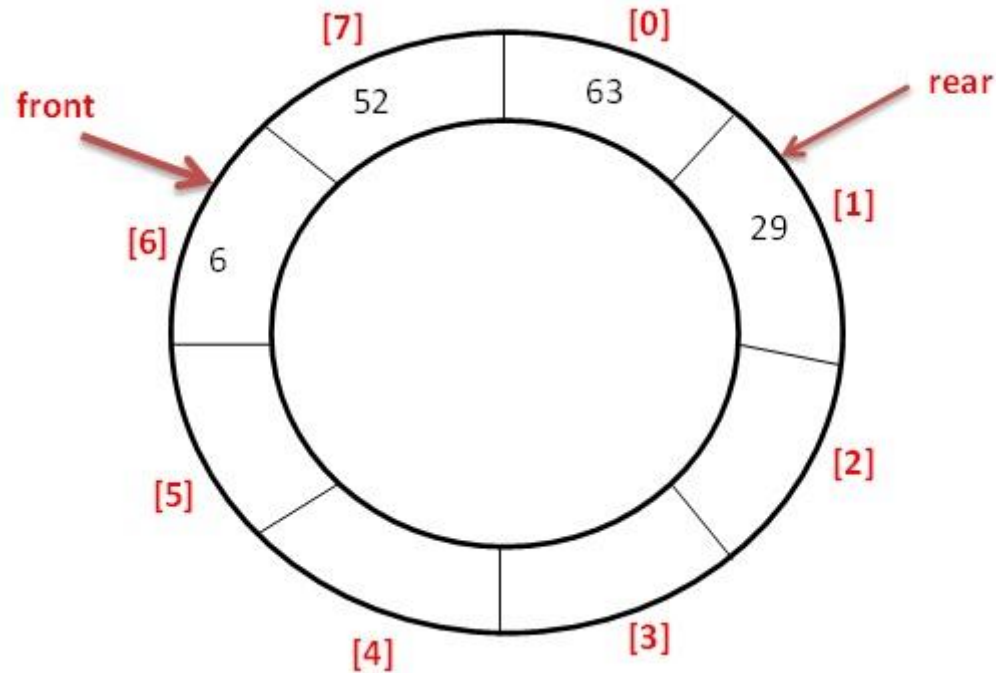
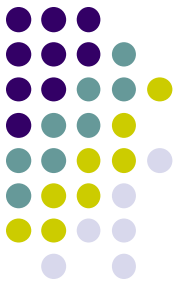
dequeue()

Circular Queues



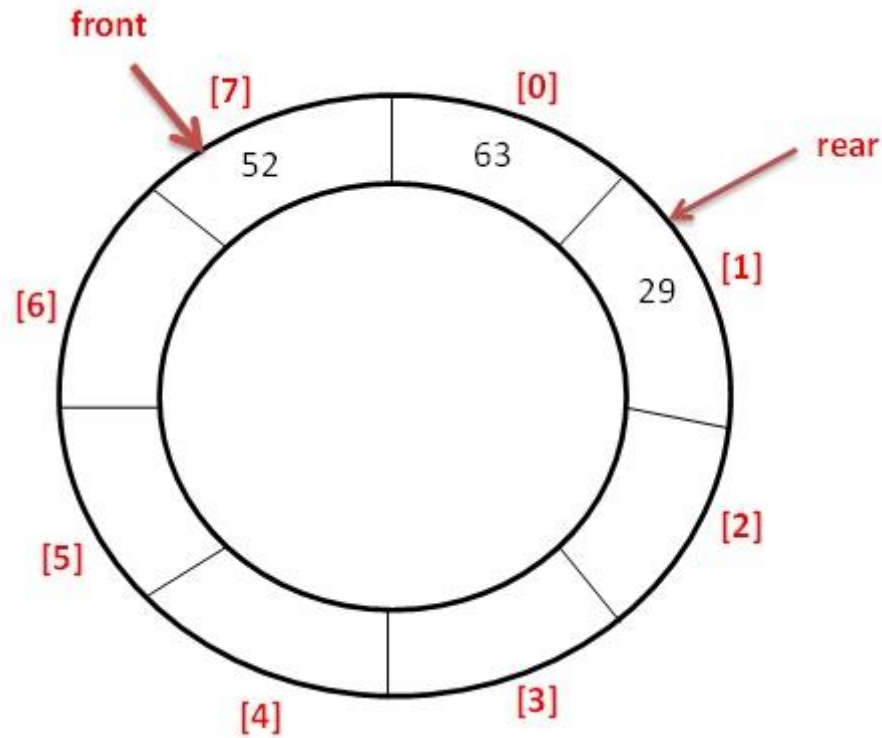
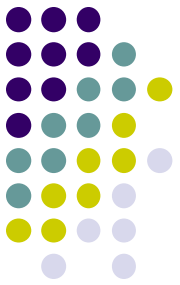
dequeue()

Circular Queues



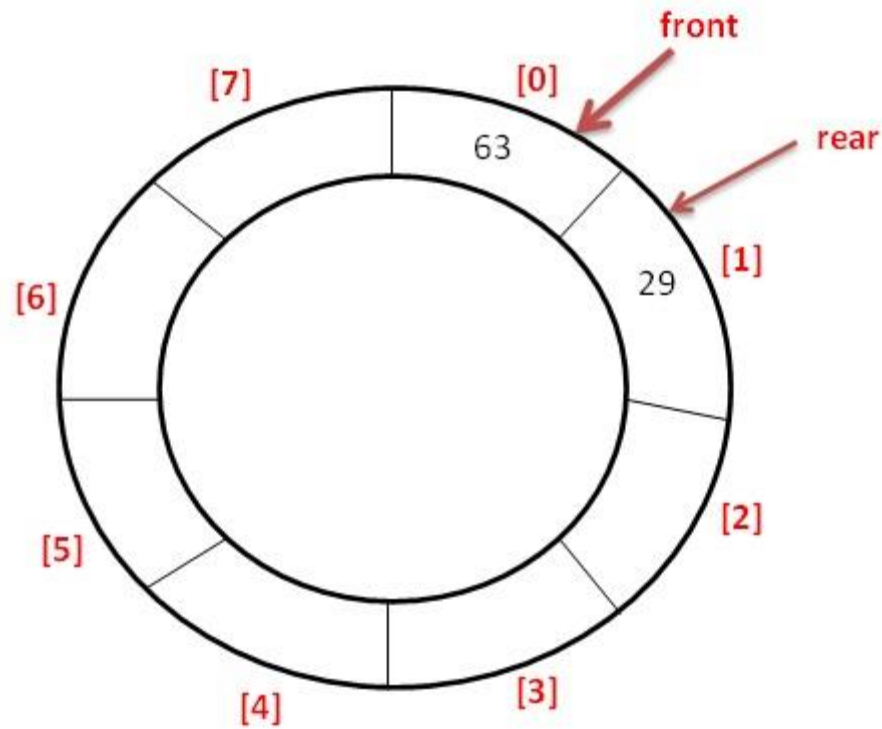
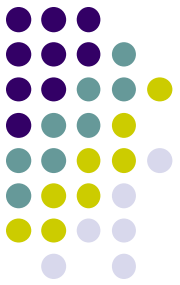
dequeue()

Circular Queues



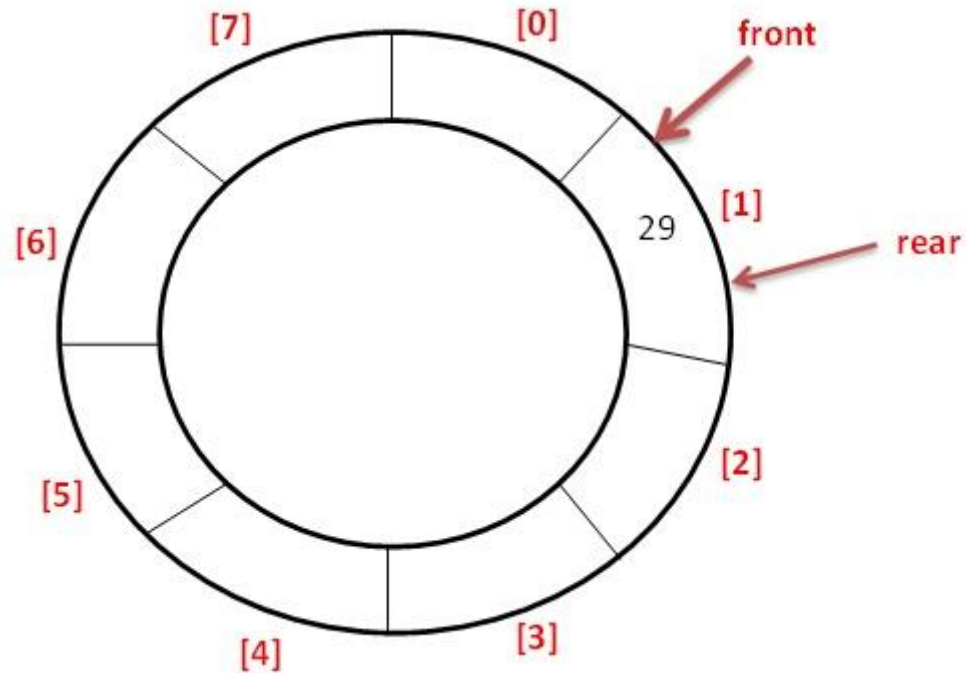
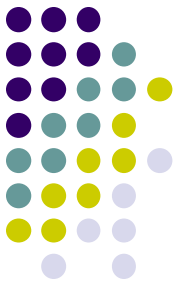
dequeue()

Circular Queues



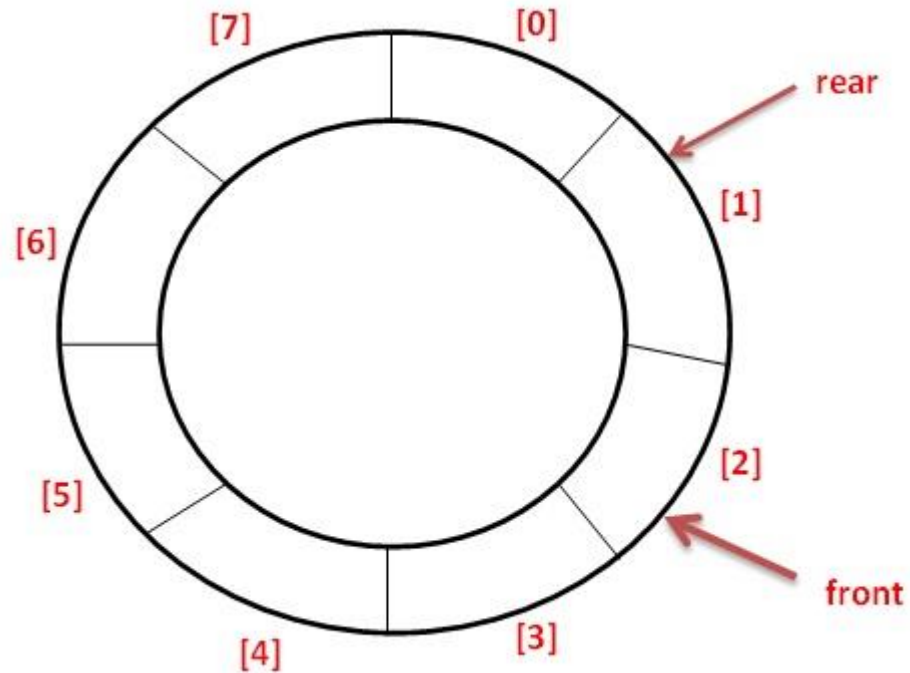
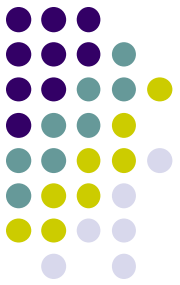
dequeue()

Circular Queues



dequeue()

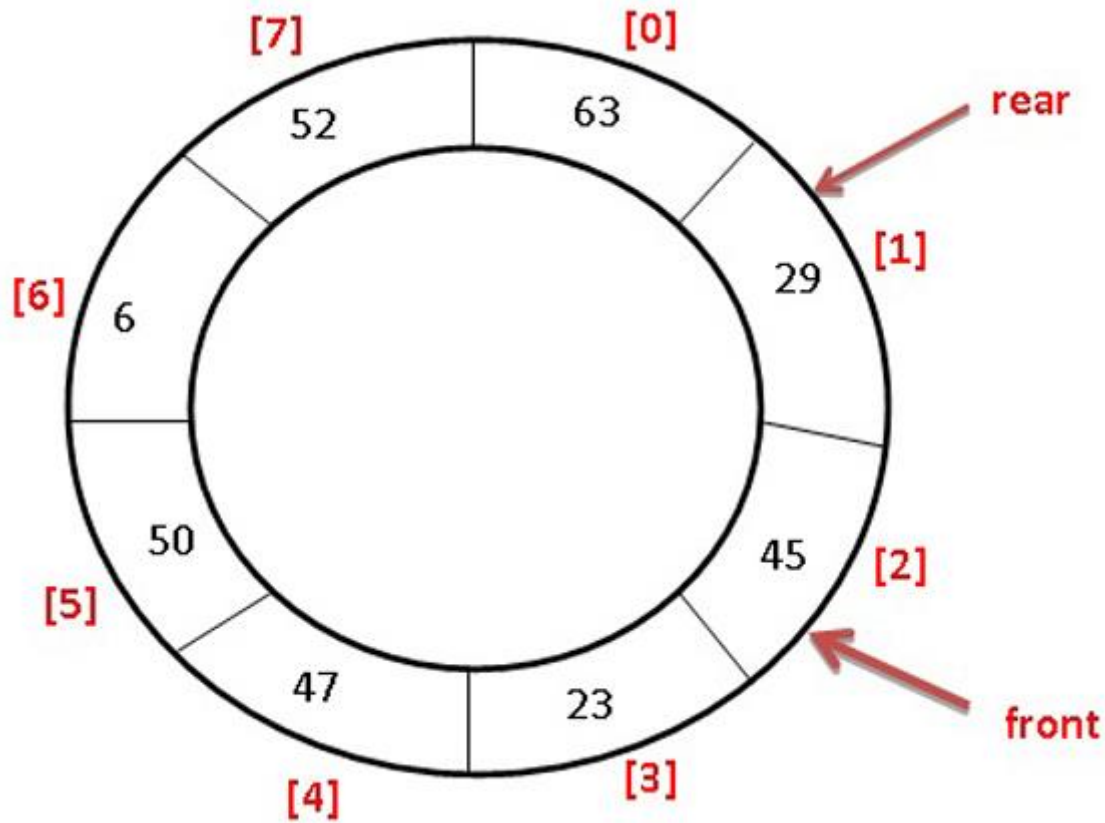
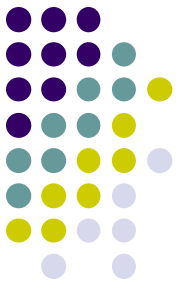
Circular Queues



Empty Queue

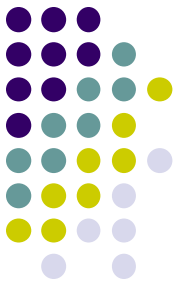
`dequeue()`

Circular Queues



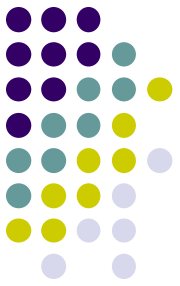
Full Queue

Array Implementation of the Circular Queues



```
typedef char QueueItemType;
class Queue{
public:
    Queue(int size);
    ~Queue();
    bool isEmpty();
    bool isFull();
    bool enqueue(QueueItemType newItem);
    bool dequeue(QueueItemType *QueueTop);
private:
    QueueItemType *items;
    int front, rear, count;
    int MaxQueue;
};
```

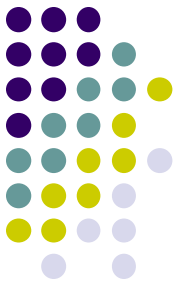

Array Implementation of the Circular Queues



```
Queue::Queue(int size) {  
    items = new QueueItemtype[size];  
    MaxQueue = size;  
    front = 0;  
    rear = -1;  
    count = 0;  
}
```

```
Queue::~~Queue(){  
    delete [] items;  
    items = NULL;  
}
```

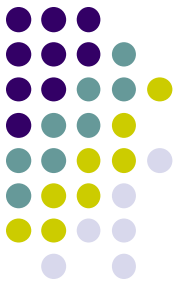
Array Implementation of the Circular Queues



```
bool Queue::isEmpty() {  
    return count <= 0;  
}
```

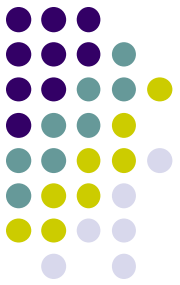
```
bool Queue::isFull() {  
    return count >= MaxQueue-1;  
}
```

Array Implementation of the Circular Queues



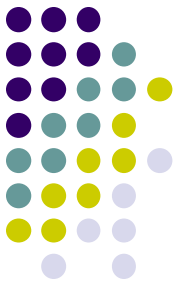
```
bool Queue::enqueue(QueueItemType newItem){
    if (isFull())
        return false;
    else{
        rear = (rear + 1) % MaxQueue;
        items[rear] = newItem;
        count++;
        return true;
    }
}
```

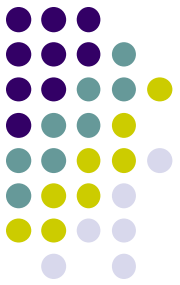
Array Implementation of the Circular Queues



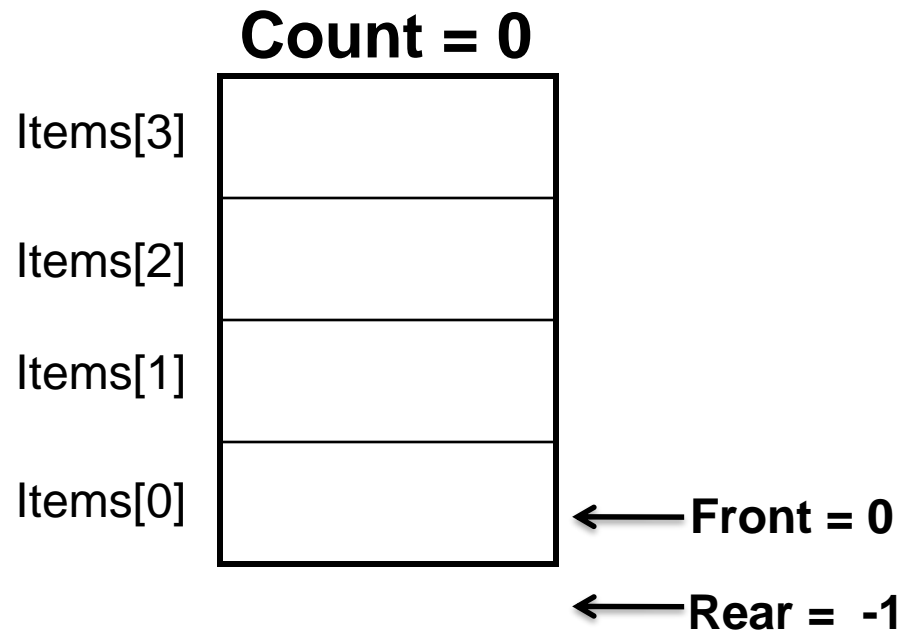
```
bool Queue::dequeue (QueueItem Type
*QueueTop){
    if (isEmpty())
        return false;
    else {
        count--;
        *QueueTop = items[front];
        front = (front + 1) % MaxQueue;
        return true;
    }
}
```

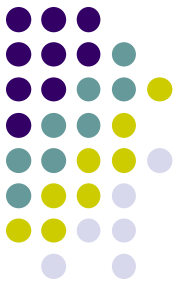
```
main(){
    QueueItem c;
    Queue Queue(4);
    Queue.enqueue('a');
    Queue.enqueue('b');
    Queue.enqueue('c');
    Queue.dequeue(&c);
    printf("%c\n",c);
    Queue.dequeue(&c);
    printf("%c\n",c);
    Queue.enqueue('d');
    Queue.dequeue(&c);
    Queue.enqueue('e');
    printf("%c\n",c);
    Queue.dequeue(&c);
    printf("%c\n",c);
    Queue.dequeue(&c);
    printf("%c\n",c);
}
```



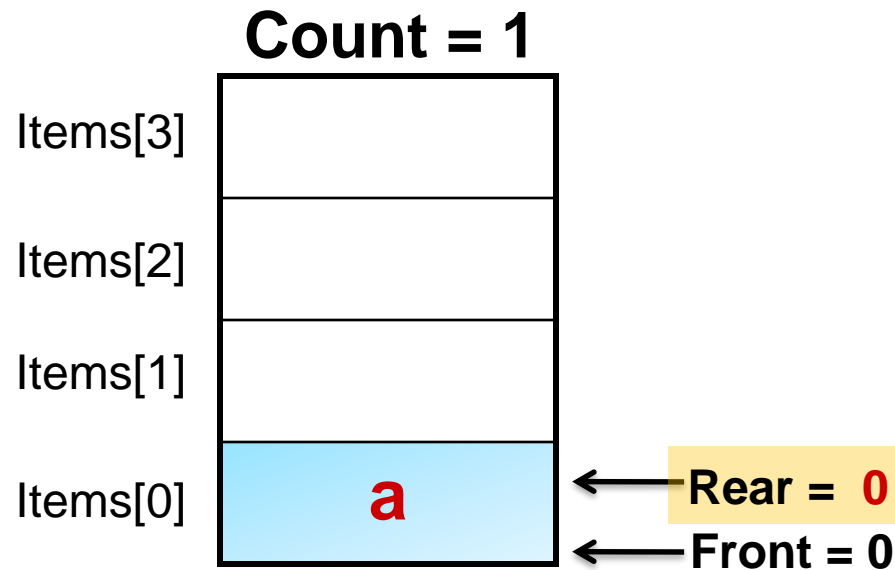


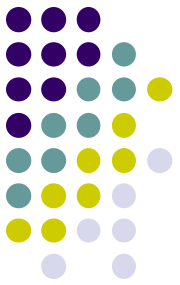
queue(4)



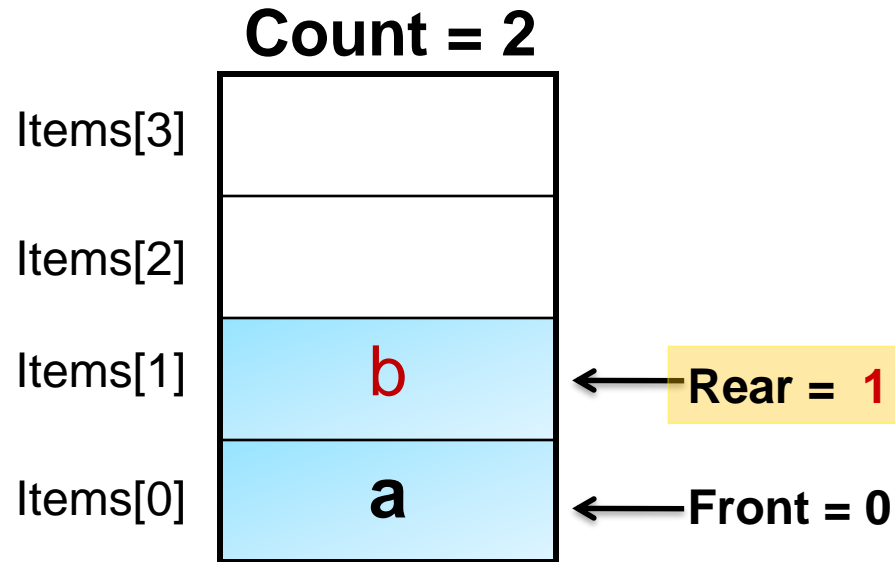


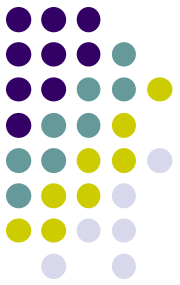
enqueue ('a')



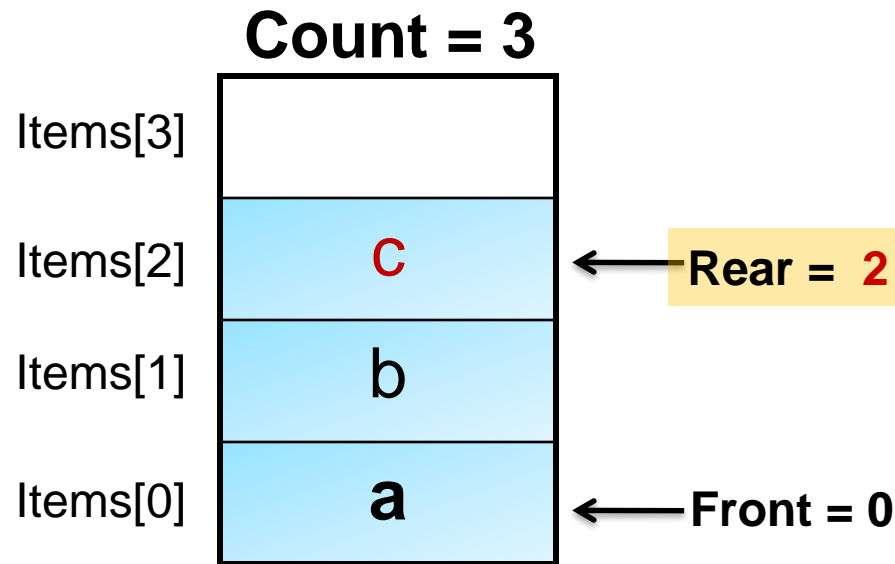


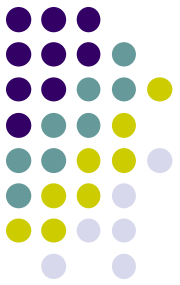
enqueue ('b')



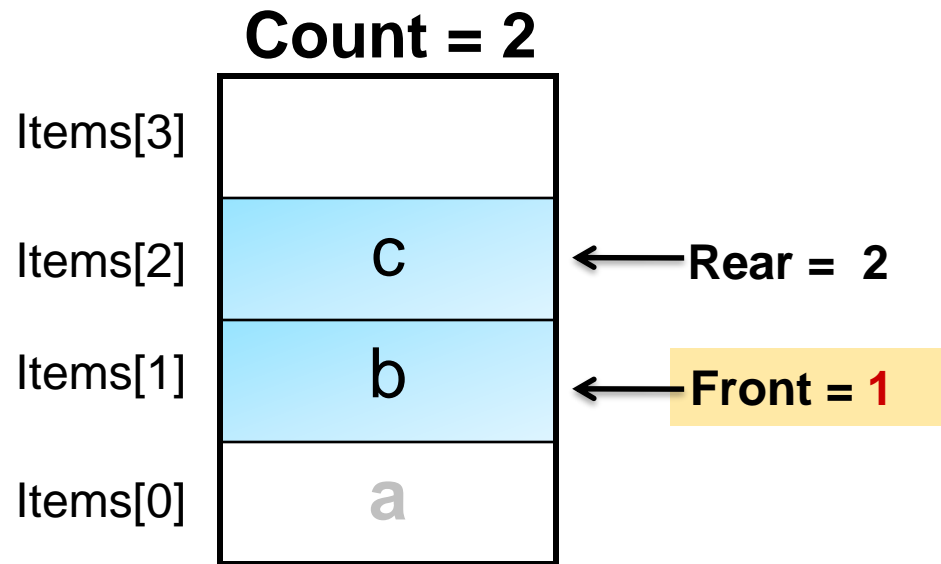


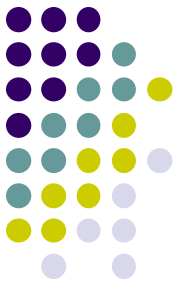
enqueue ('c')



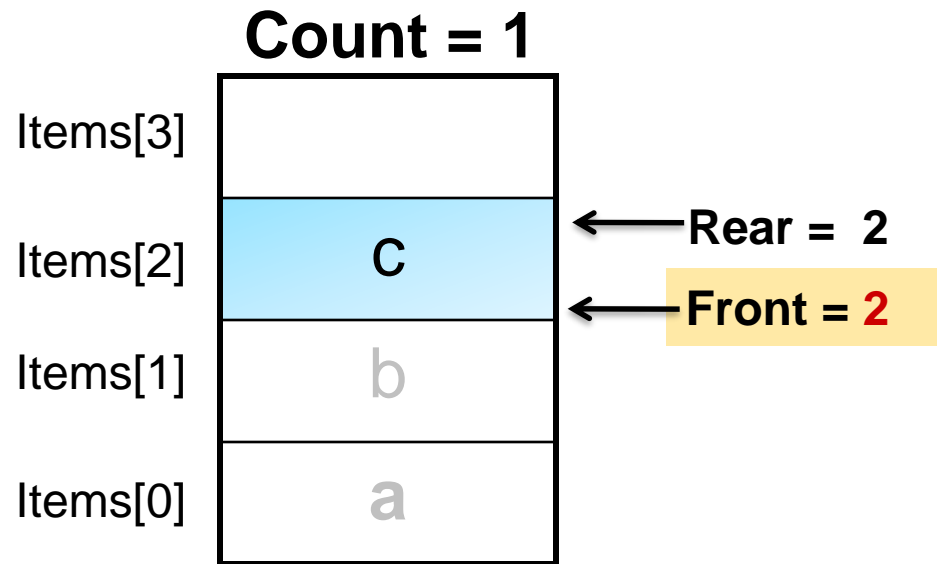


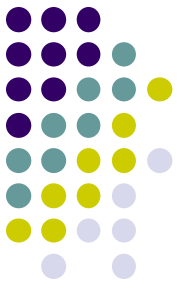
dequeue (&c)



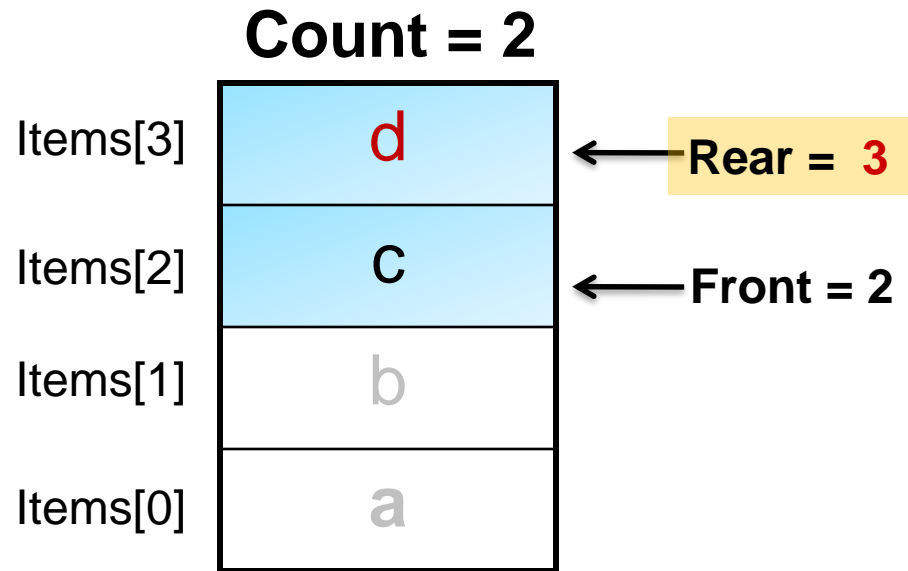


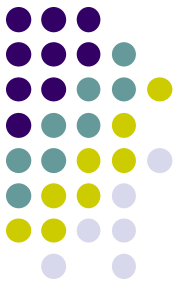
dequeue (&c)



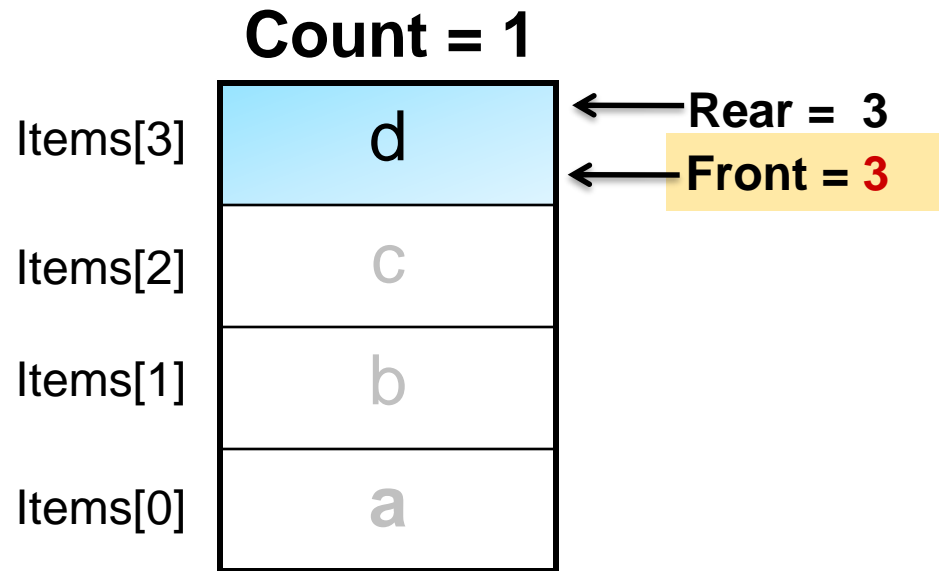


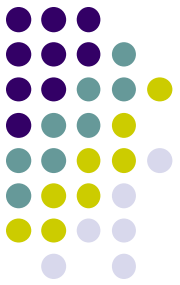
enqueue ('d')



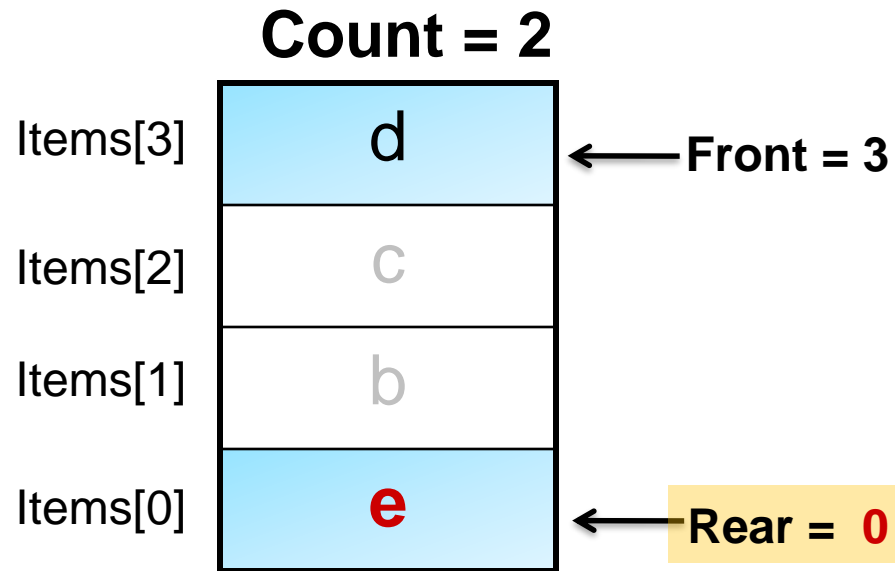


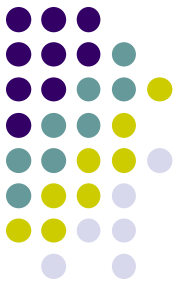
dequeue (&c)



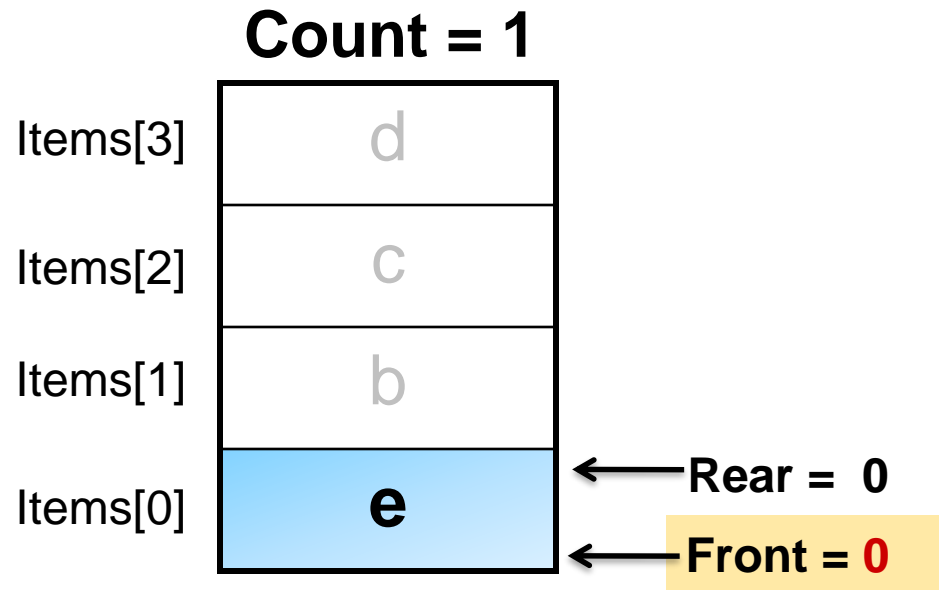


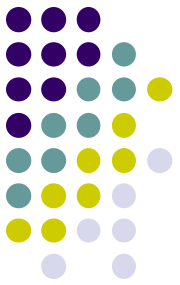
enqueue ('e')



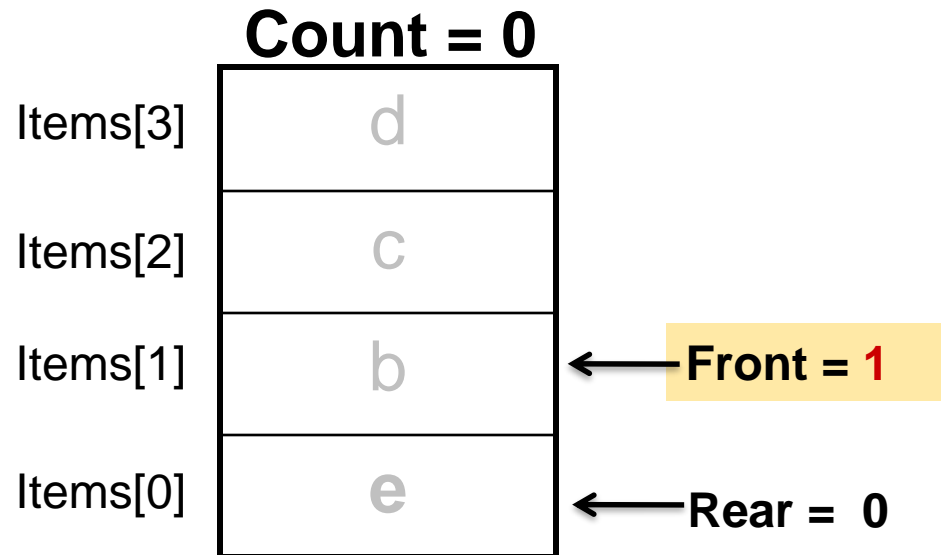


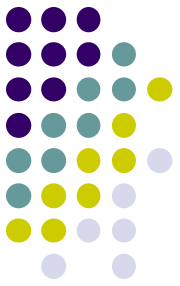
dequeue (&c)





dequeue (&c)

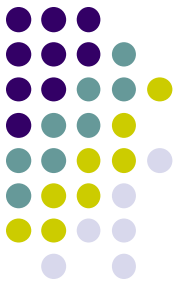




Applications of Queues

- Queues are used in operating systems, for controlling access to shared system resources such as:
 - Printer
 - Disk access on a network
 - CPU

Printer Queue

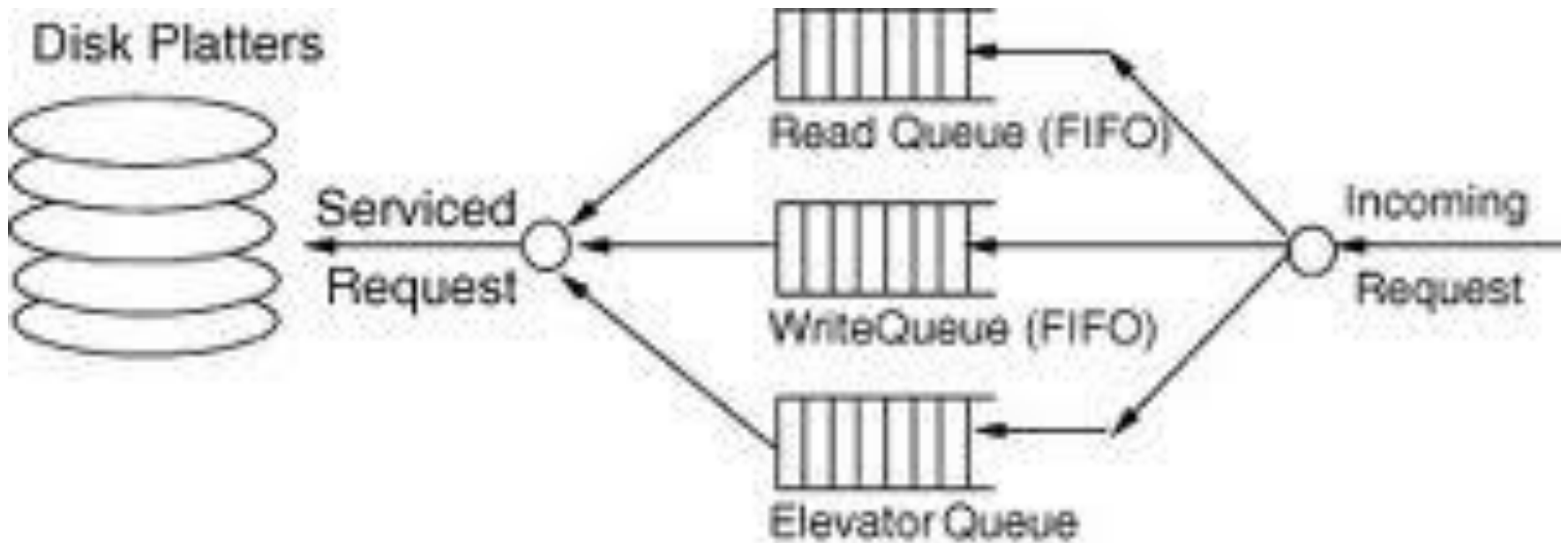
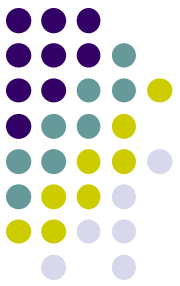


HP Color LaserJet CP1510 series PCL6 - Use Printer Offline

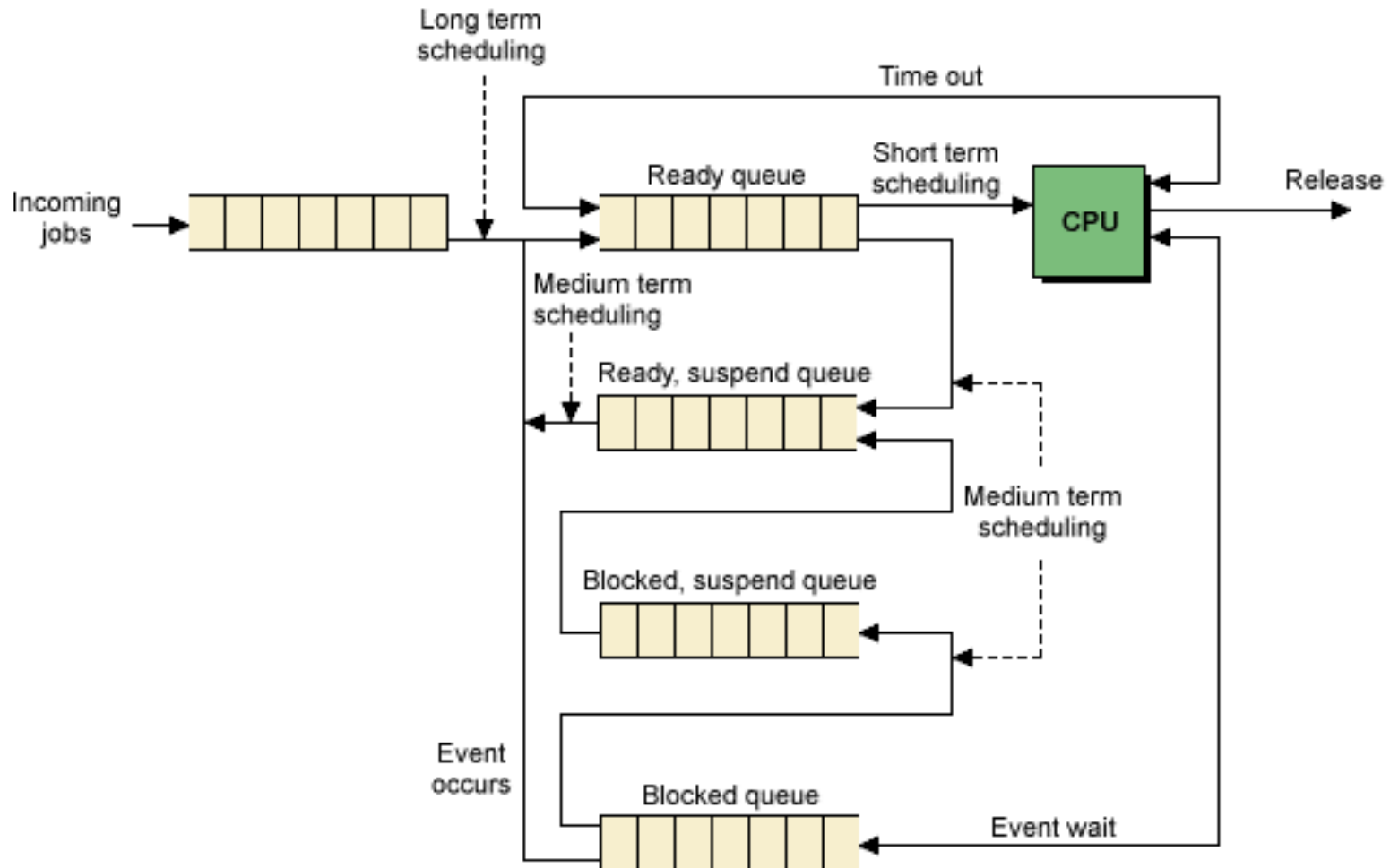
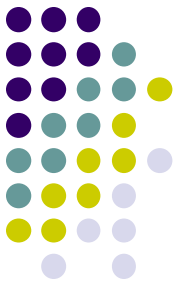
Document Name	Status	Owner	Pages	Size	Submitted	Port
Microsoft Word - prior quese.doc		Dr G S Lehal	18	6.54 MB	09:09:54 11-06-2012	
Microsoft Word - Queue Script....		Dr G S Lehal	20	2.60 MB	09:09:08 11-06-2012	

2 document(s) in queue

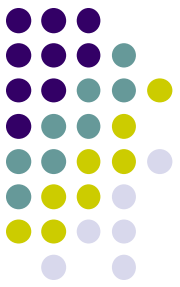
Disk Queue



Process Queues



Queues in Simulation



teller 1



teller 2

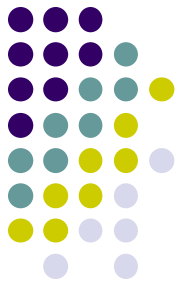


teller 3

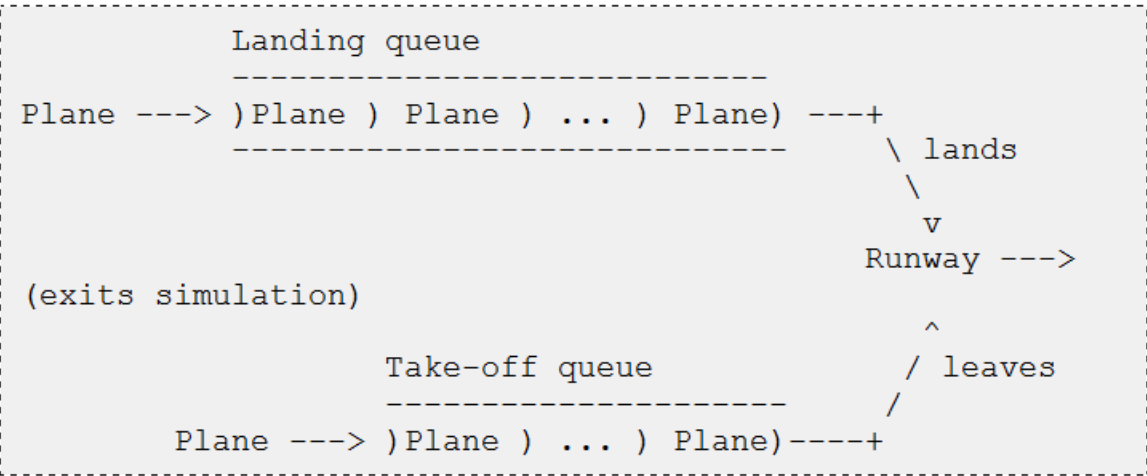


teller 4

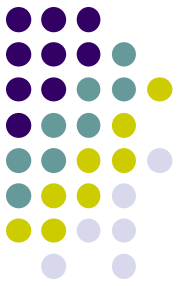




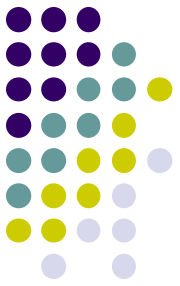
Simulation of Airplane Traffic



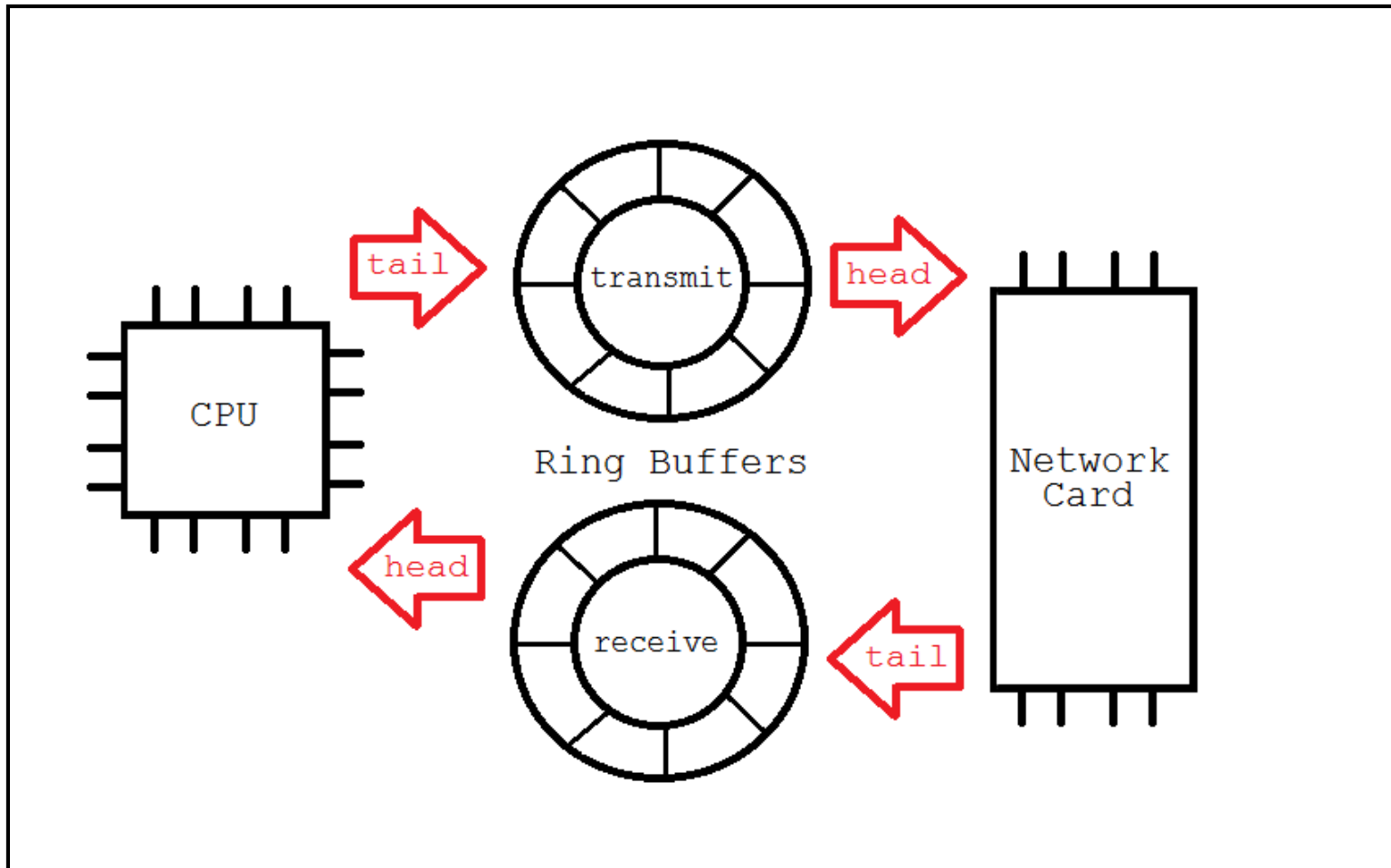
Other Applications of Queues

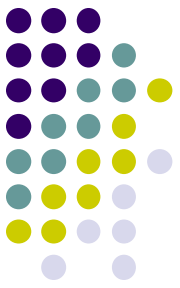


- Reading from the keyboard
- Implementing buffer
- Level order traversal of trees or breadth first search in a graph
- Implementing Radix Sort

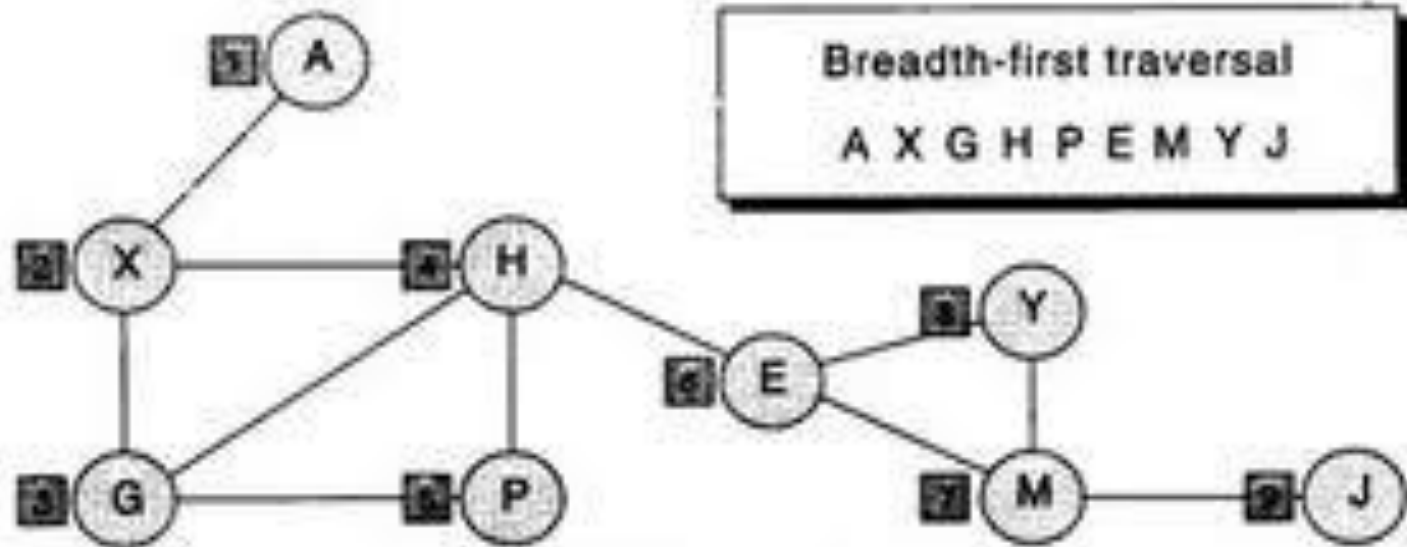


Buffers





Breadth First Traversal of Graph



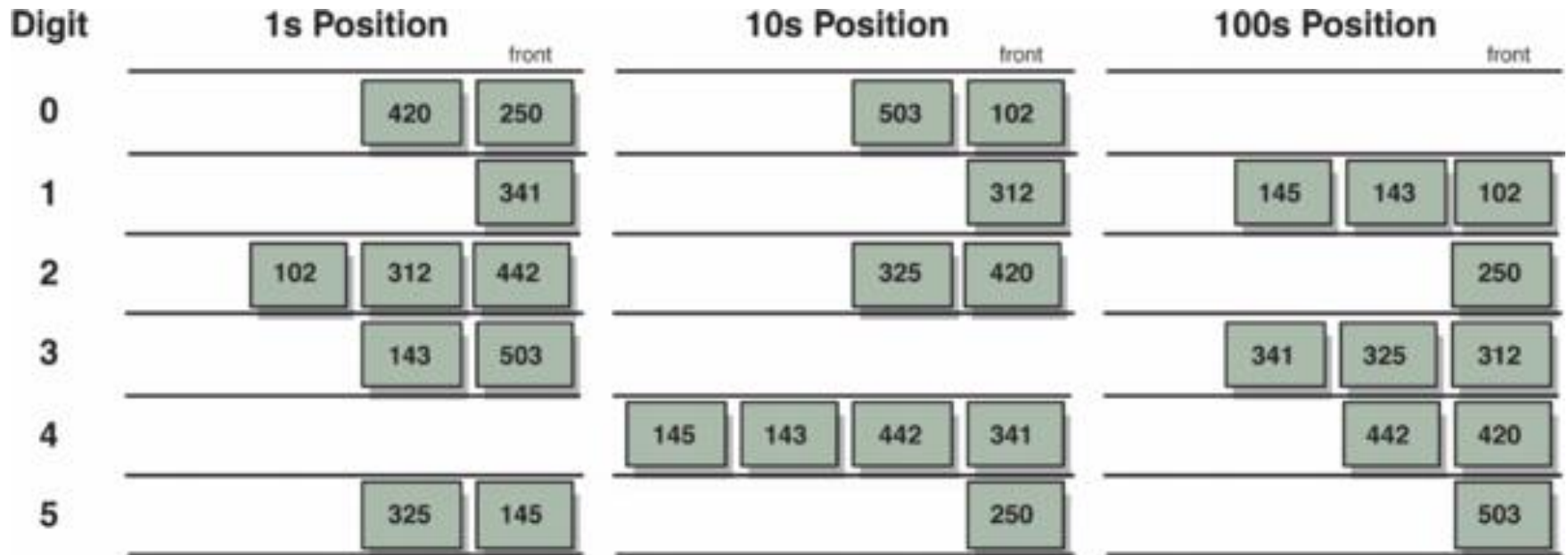
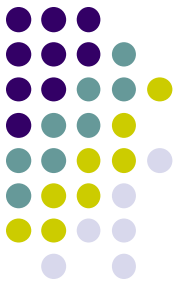
Breadth-first traversal
A X G H P E M Y J

(a) The graph



(b) Queue contents

Radix Sort



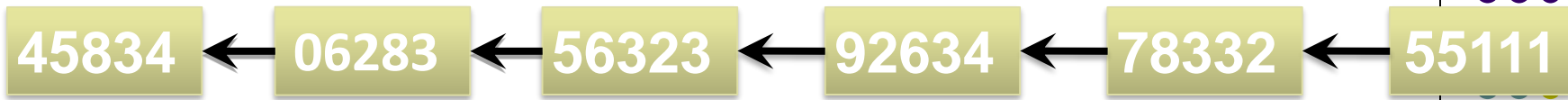
Original List:



Radix Sort Algorithm



1. Place all the integers in the main queue.
2. Remove each value in the main queue and place it in a digit queue corresponding to the digit being considered, starting with the least significant digit.
3. Once all the values are placed in the appropriate digit queue, collect the values from queue 0 to queue 9, and place them back in the main queue.
4. Repeat steps 2 and 3 with the tens digit, the hundreds digit, and so on. After the last digit is processed, the main queue contains the values in ascending order.



0
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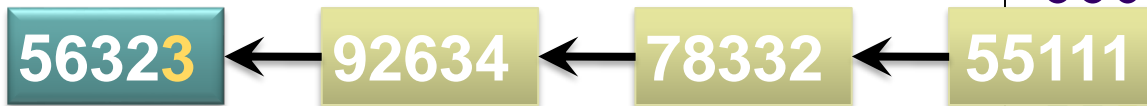


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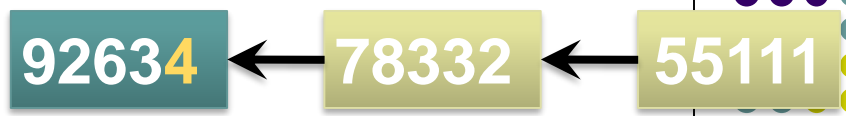
45834



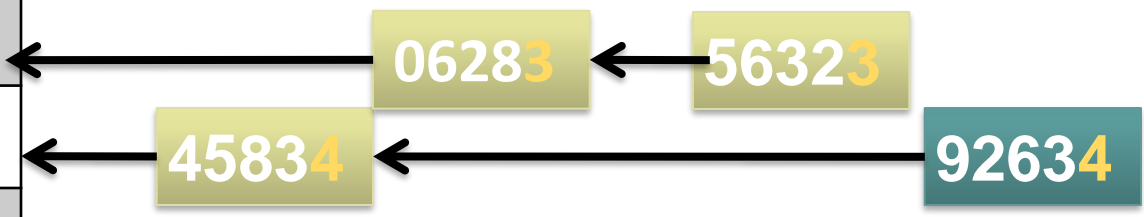


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78332 ← 55111

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78332

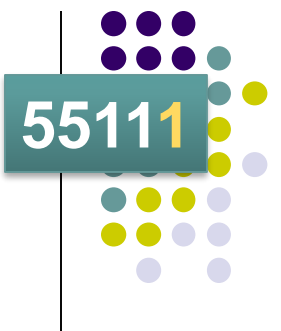


06283 ← 56323



45834 ← 92634





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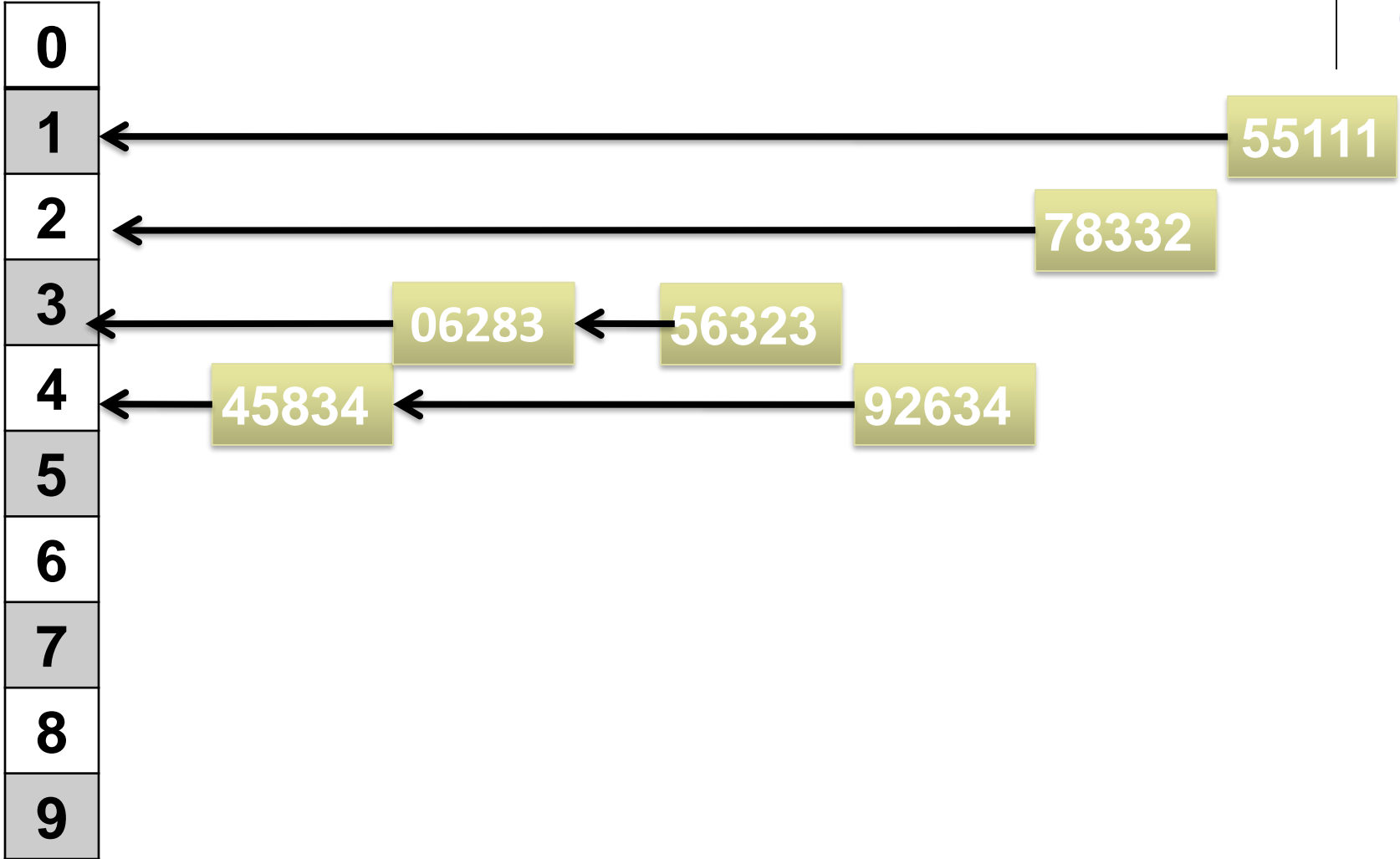
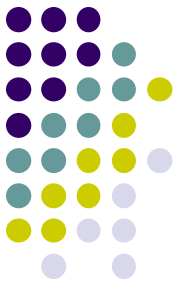
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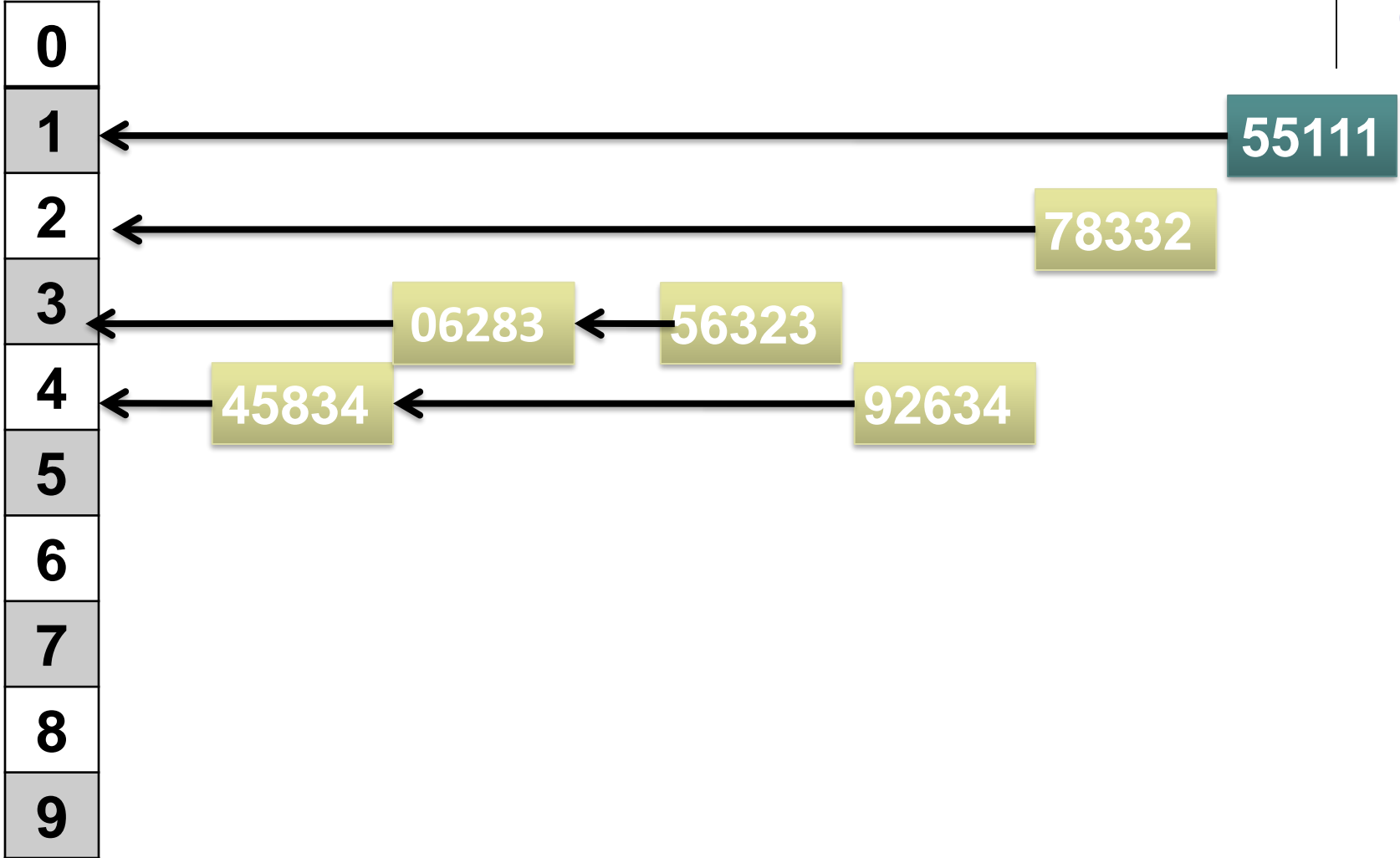
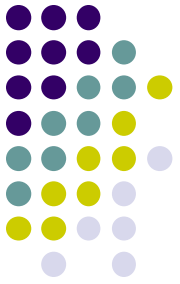
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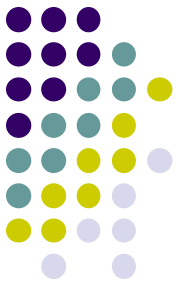
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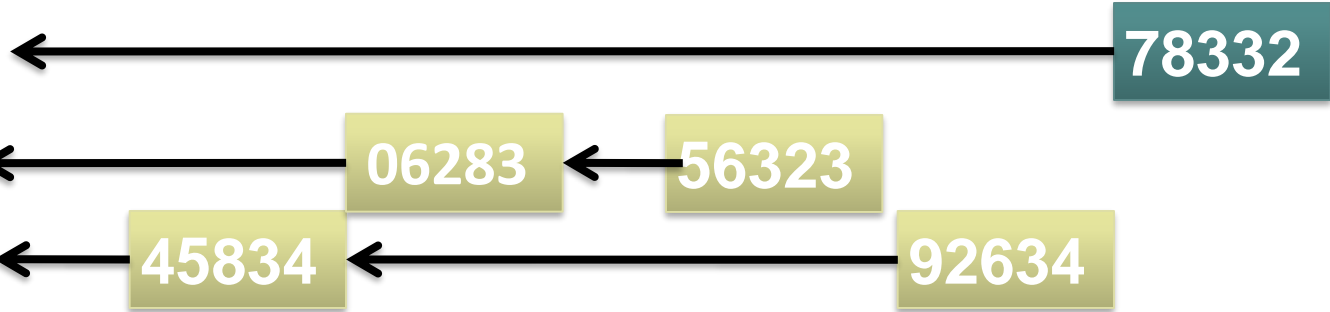
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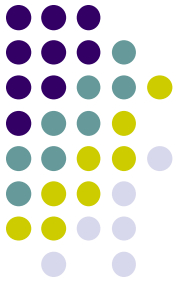
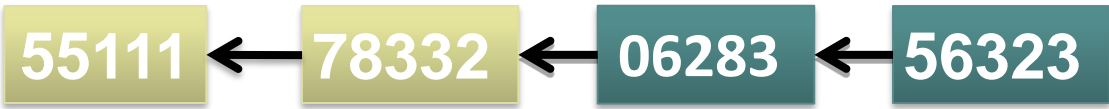




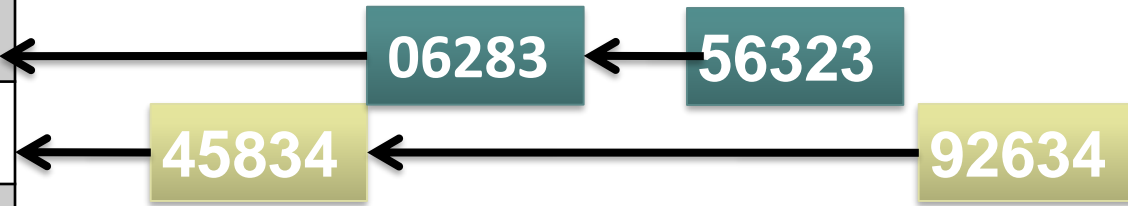
55111 ← 78332

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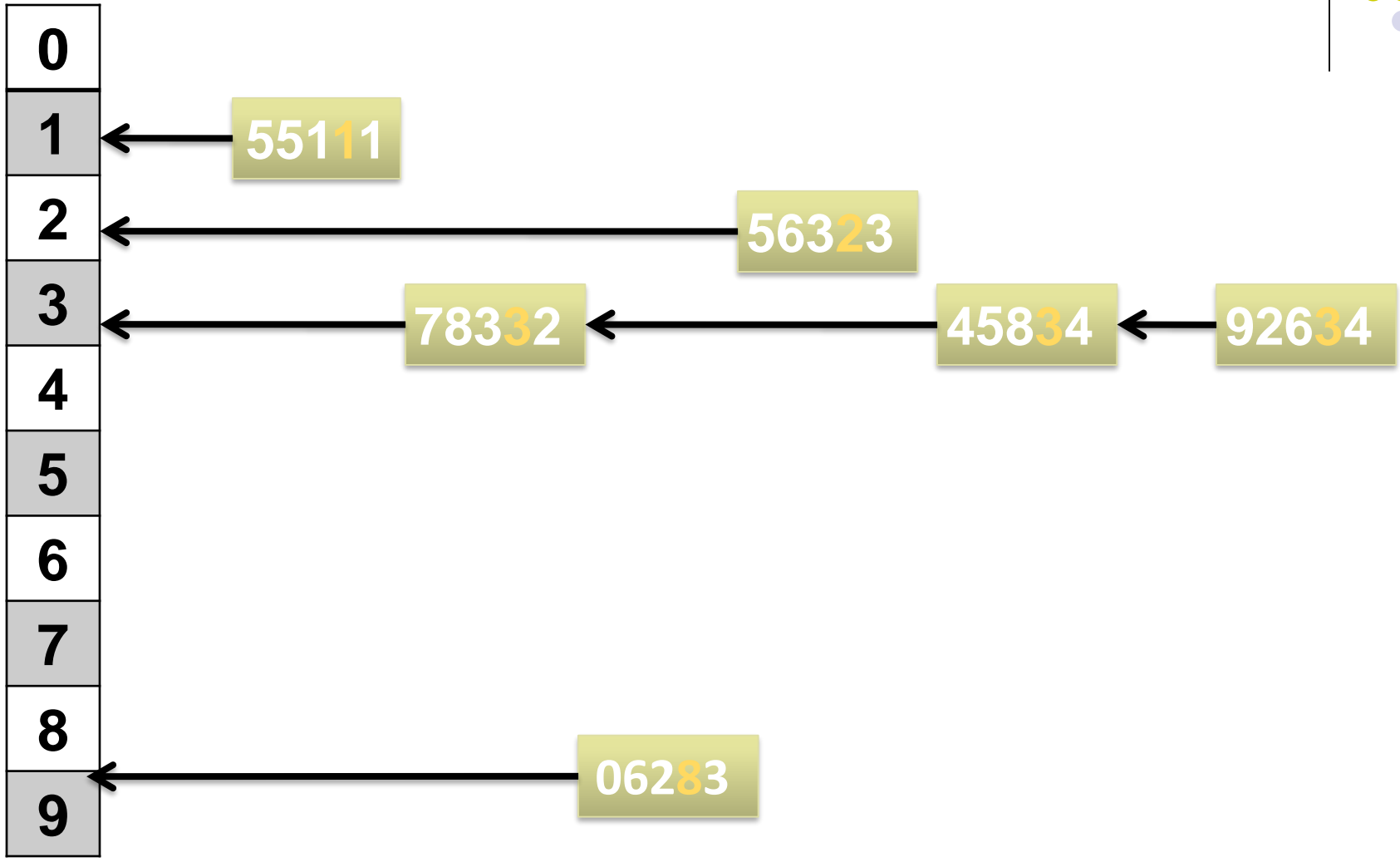
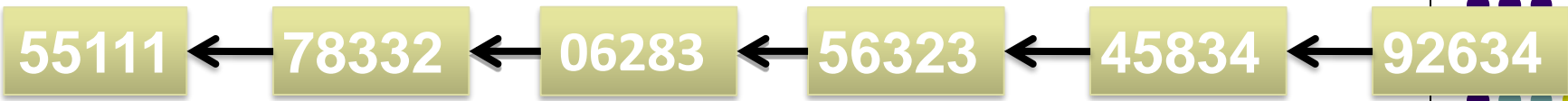
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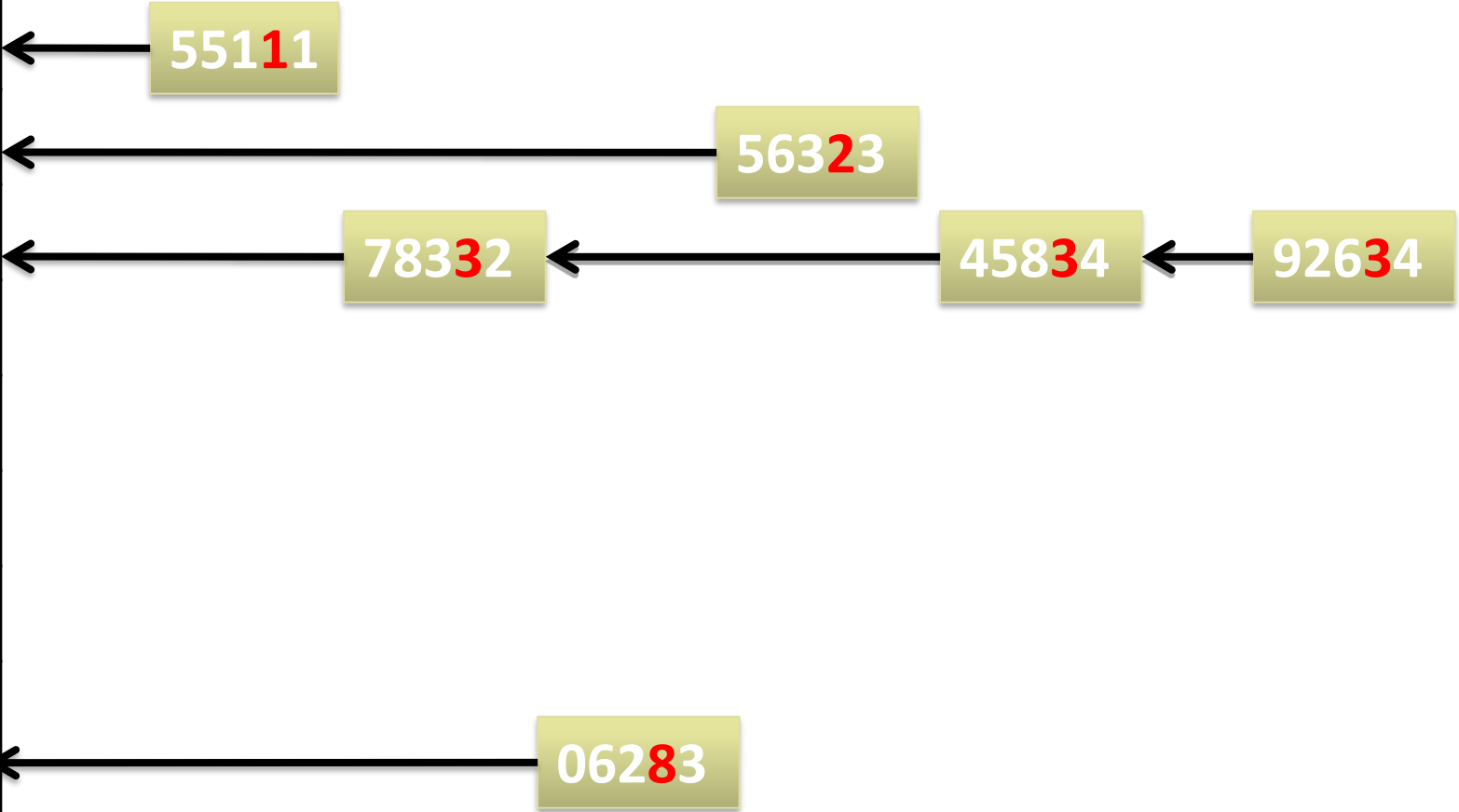
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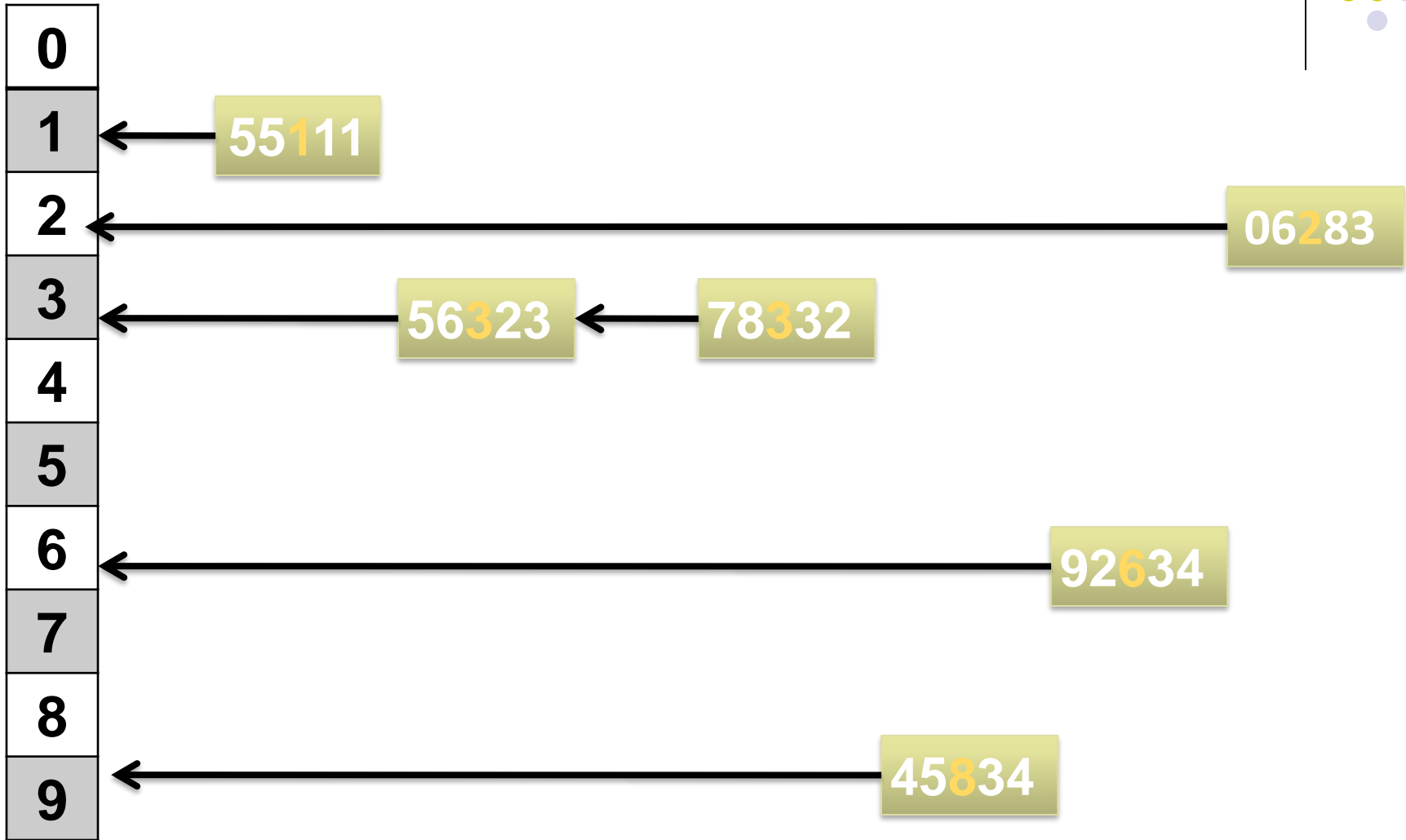
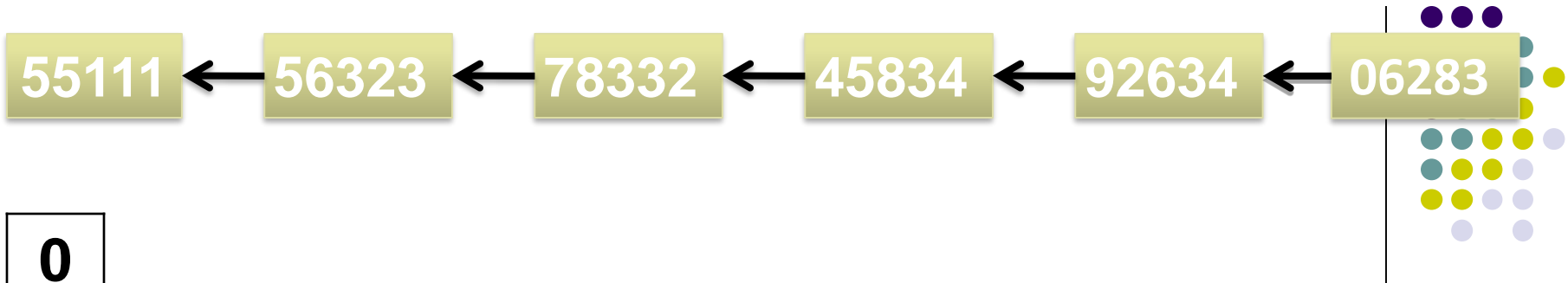






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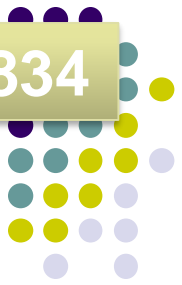






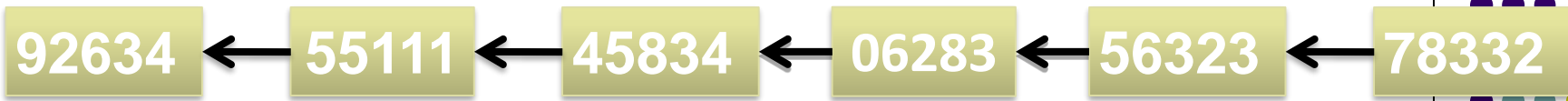
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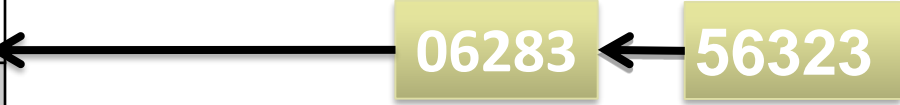


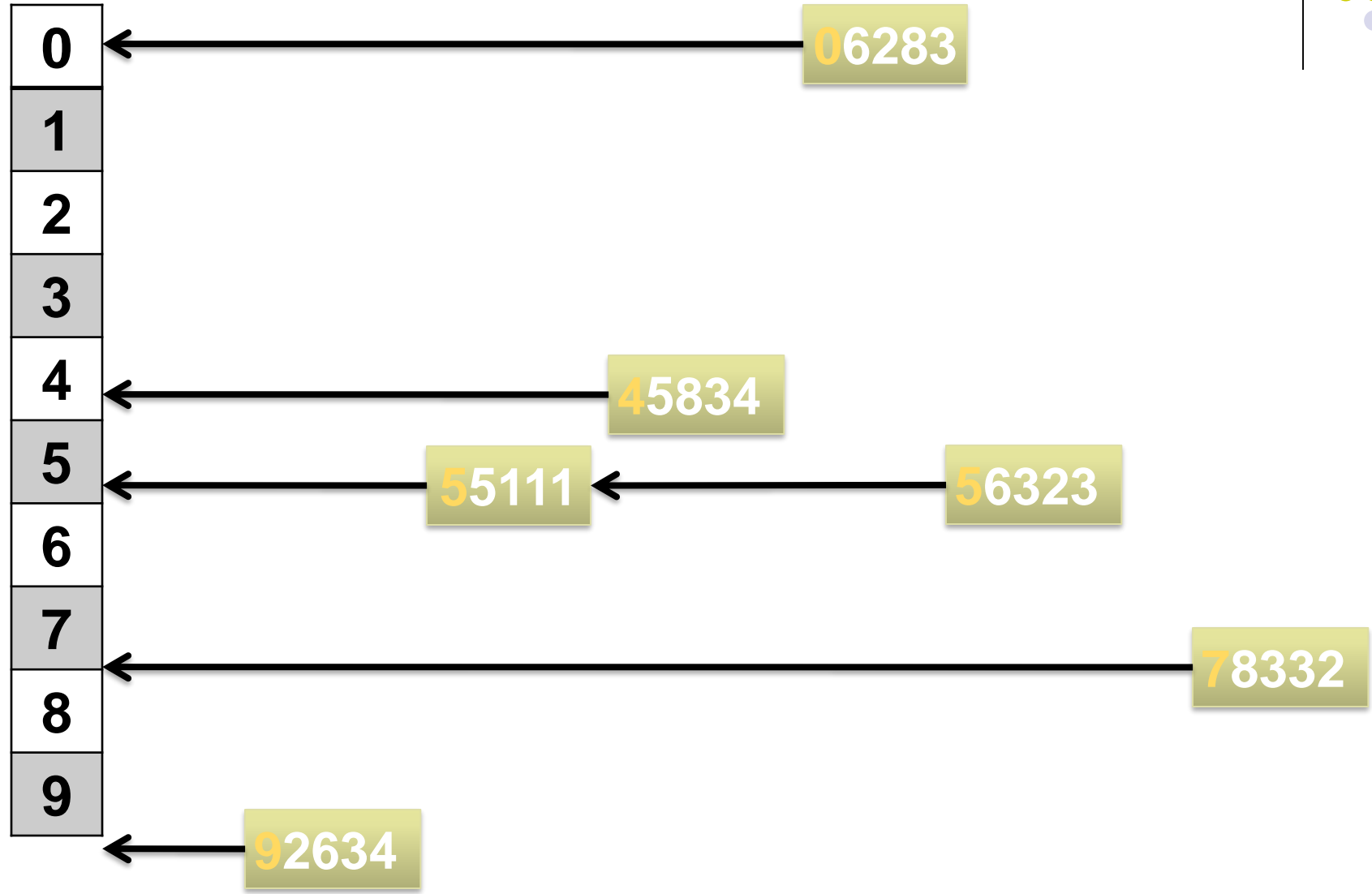
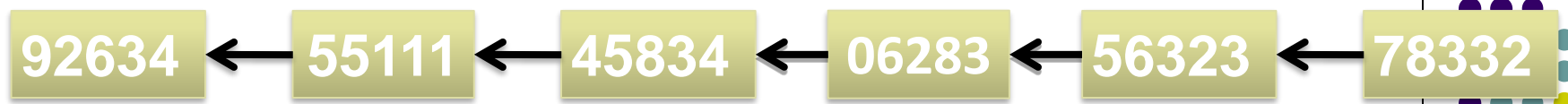
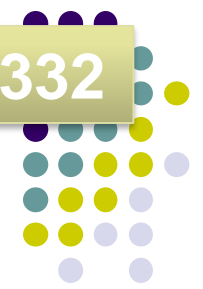
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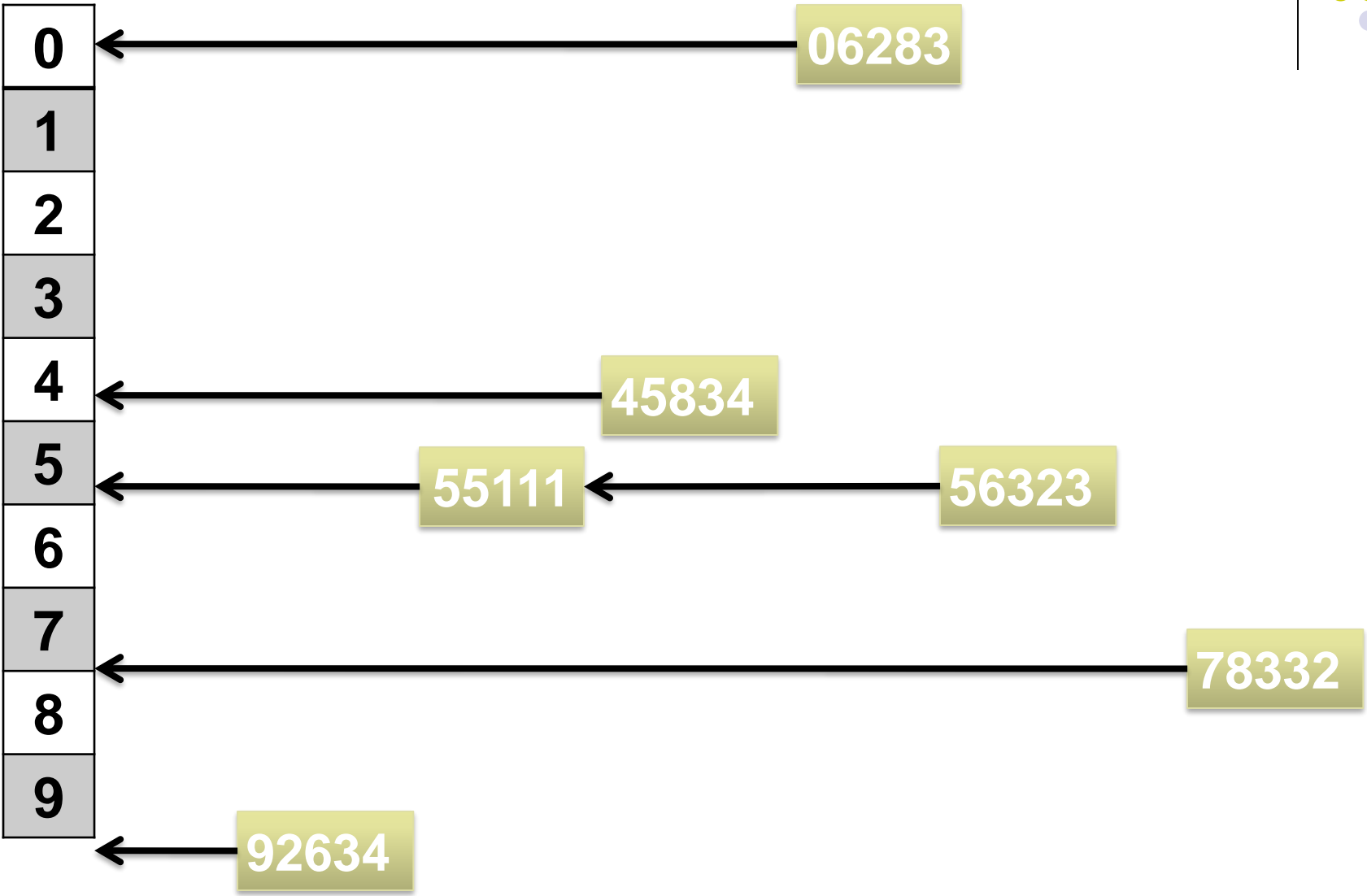




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