

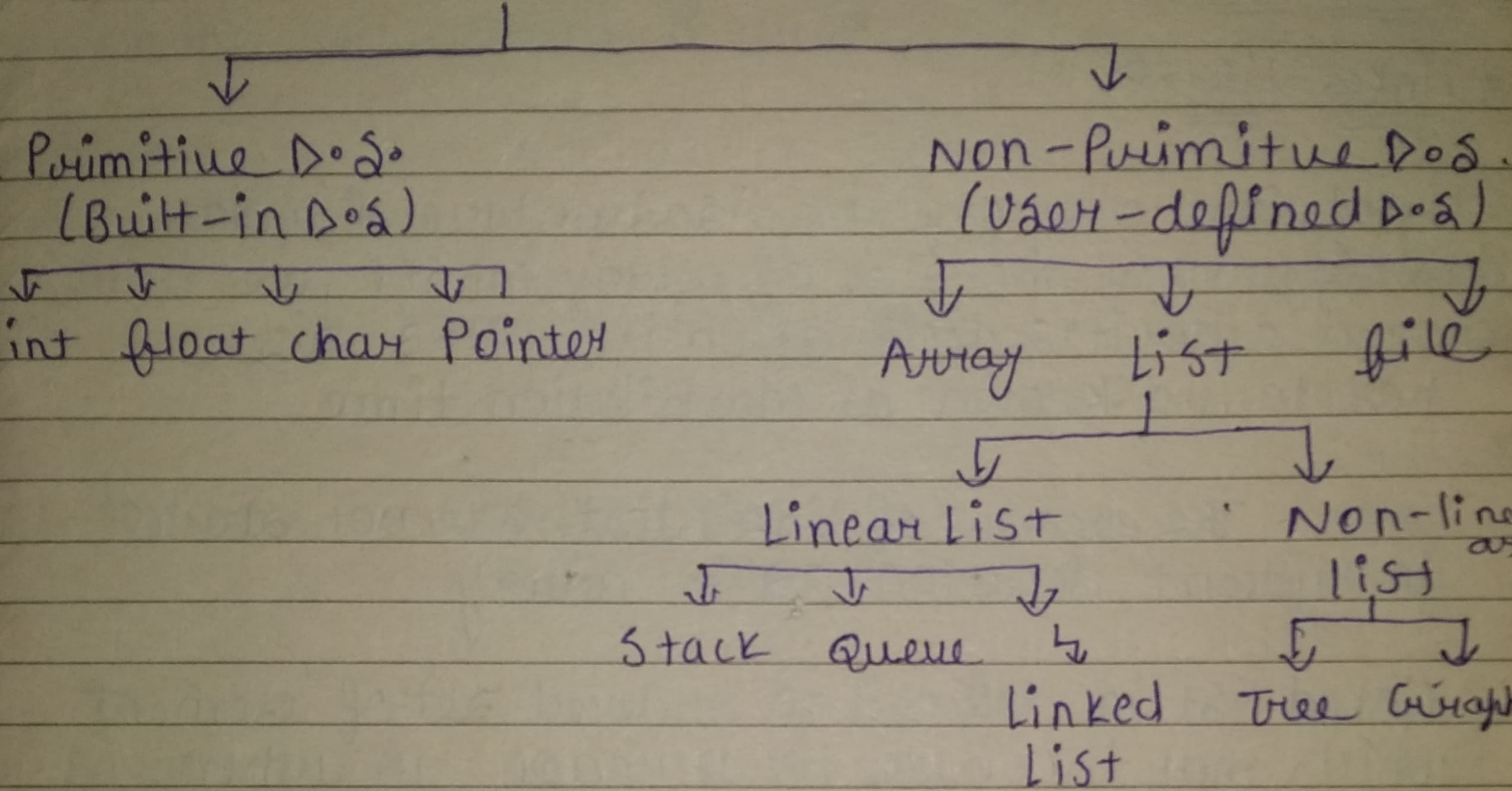
Data Structures

Data may be organized in many different ways. The logical or mathematical model of a particular organization of the data is called data structure.

Operations \Rightarrow :

1. Insert \Rightarrow Adding a new record to D.S.
2. Delete \Rightarrow Removing a Record from the D.S.
3. Searching \Rightarrow Finding the location of the record with a given key value or finding the location of all records which satisfy one or more condition.
4. Traversing \Rightarrow Excessing each record exactly once ~~so~~ so that certain items in the record may be proceed.
5. Sorting \Rightarrow Arrange the Record in some logical order.
combine
6. Merge \Rightarrow Combine the data of two different sorted files into a single sorted file.

Data Structure



1. Primitive D.O.S. ⇒ is directly operated by the machine instruction and the size of this data structures is pre-defined. Ex ⇒ int, float, char, string

2. Non-Primitive D.O.S. ⇒ is a user-defined data structures and these are derived from the Primitive D.O.S.

The non Primitive D.O.S. emphasize on structuring of group of homogenous & heterogenous data items.

For ex ⇒ Array, stack, queue, Linked List etc.

Date: →

Linear List →

Linked-List →

Linked list overcome the disadvantage of array data structure the size of array cannot be changed after its declaration i.e. its size has to be known at compilation time.

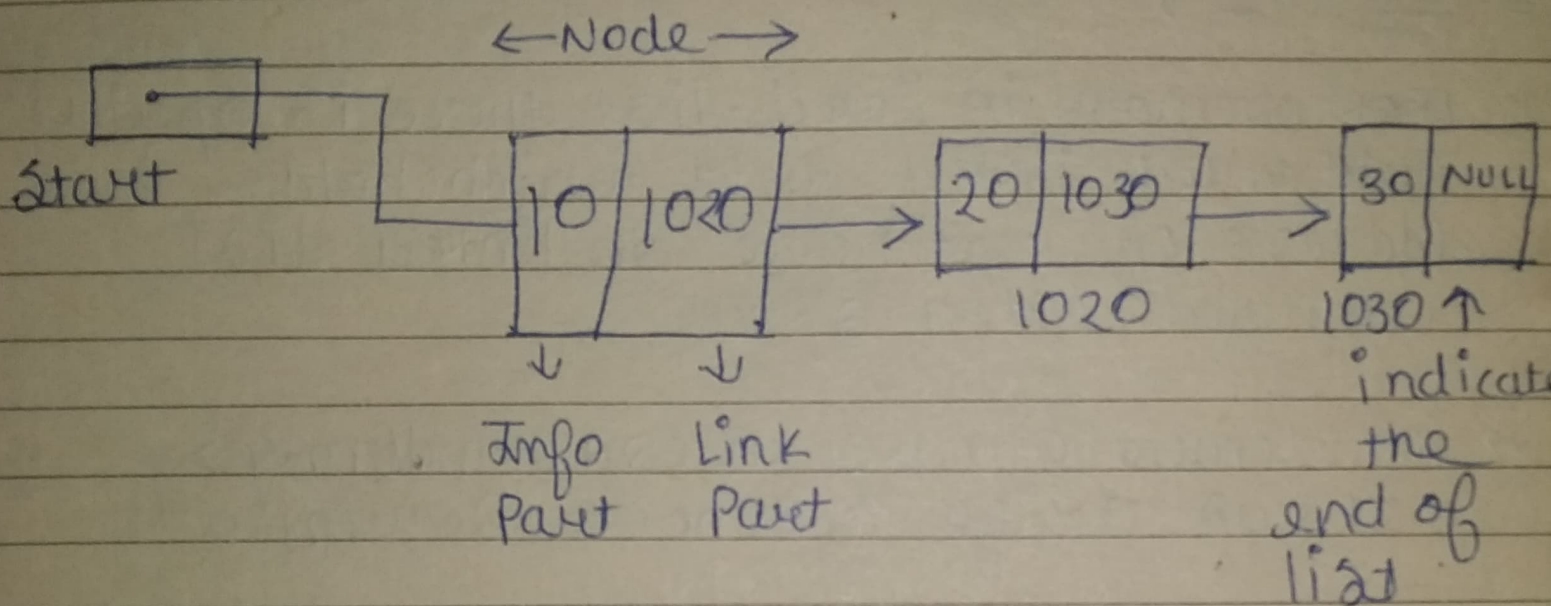
The element of linked list are not stored in adjacent location as in array.

Linked list is defined as ordered set of element which may increase or decrease as when required. In linked list memory space located for the element of the list can be extended at any time.

A linked list is an ordered collection of homogenous data elements, called nodes, where the linear order is maintained by means of links of pointer.

Each node is divided into two parts. The first part contain information and second part contain the address of next node in list.

The number of pointer is maintained depending on the requirement & usage.



Advantage :-

1. Linked list are dynamic data structure.
2. They can grow & shrink during the execution of program.
3. Efficient memory utilization → memory is not allocated & memory is allocated whenever required.
4. Insertion & deletion of node are efficient → linked list provide in insertion data item have specified posⁿ or deletion of data item for

given position.

* Disadvantage \Rightarrow :

- 1^o More memory in linked-list there is special field called linked field which hold address of next node so linked list required extra space.
- 2^o No element can be accessed randomly \Rightarrow It has access each node sequentially.
- 3^o Reverse traversing is difficult in linked list.

Application \Rightarrow

1^o Linked list is used to represent Polynomial expression.

2^o Linked list is used to symbolic table, Symbol table is data structure used in compiler for keeping the value of variable & constant data used in appl. of our program.

3^o Linked list are used to represent sparse matrix.

\rightarrow no. of elements non-zero more

L.P. → Recursive Node
R.P. → Successive Node

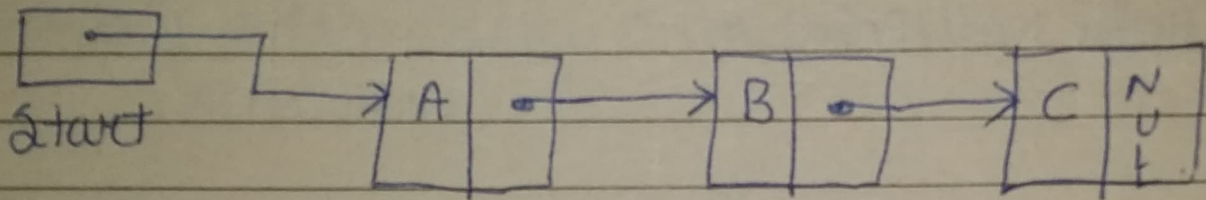
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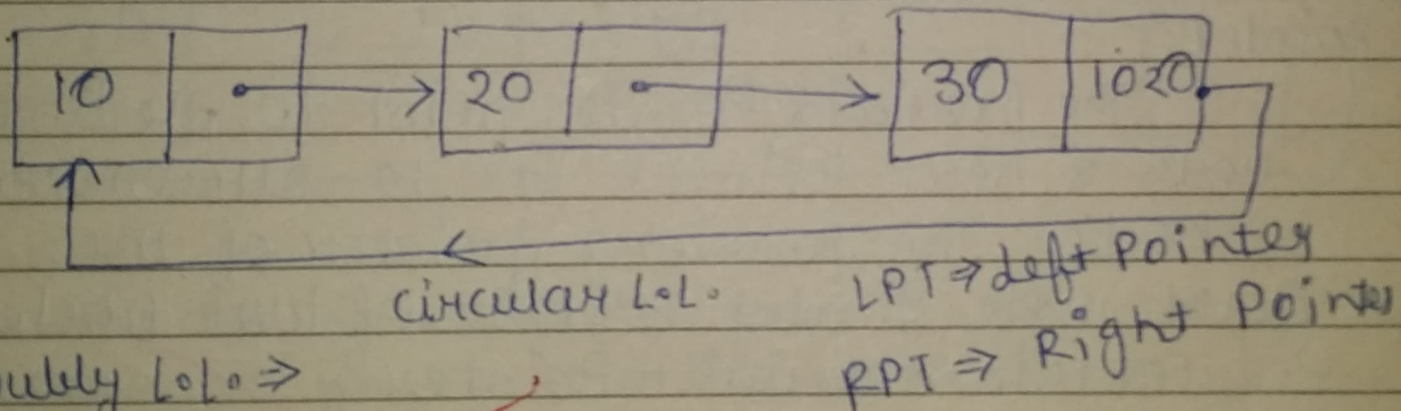
Date ⇒ 23/07/18

Types of Linked List ⇒ :

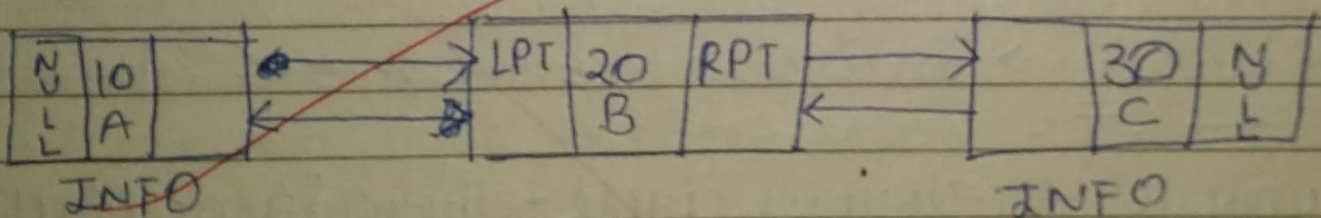
1. Singly Linked List ⇒ :



2. Circular L.L. ⇒ :



3. Doubly L.L. ⇒ :



1. Singly-Linked List \Rightarrow In this linked list each node is divided into two parts. The first part contains the information & 2nd part contains address of the next node in the list. Each node has a single pt. to next node. and the last node contains null pointer.

2. Doubly L.L. \Rightarrow are also called to bay list. In this l.l. each node has two pointers previous & next pointer.

The previous pointer point to predecessor node & next pointer pt. to successor node. The previous pointer of the first node & next pointer of last node contain the null-pointer.

3. Circular L.L. \Rightarrow In circular L.L. ~~the~~ ~~previ~~ all nodes are connected to form a circle. There is no Null pointer at the end of the list the elements point to each other in a circular way which form a circular chain & the last node contains the address of the 1st node.

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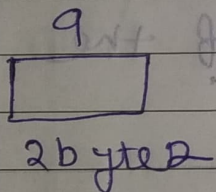
Introduction to linear Data Structures:

why array concept was introduced:

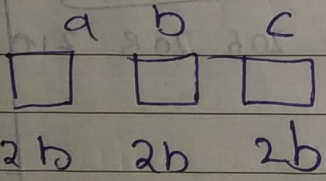
Example ->

we create a int variable in c++

```
int a;
```



```
int a, b, c;
```



For reduce time the array concept was introduced

```
int a[100]
```

if we want to make same

Array ->

Array is defined as

a collection of similar type elements.

This means an array can store either all integers & all floating point, all characters, but all of same data type.

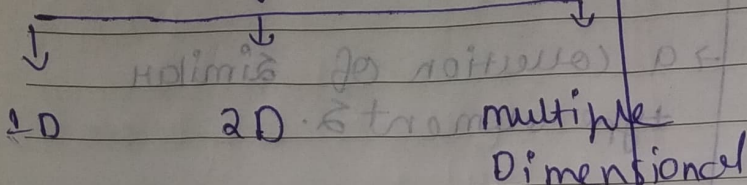
Some points about array:

Array are always stored in consecutive memory locations

Array name is actually a pointer to the first location of the memory block allocated to the name of array

Types: →

Array



```

int a[5]; int b[2][5]; int c[2][3][3];
  
```

one-dimensional array

→ only one subscript specification is needed to specify

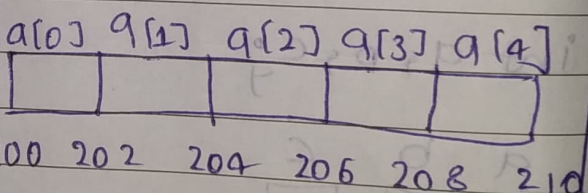
→ declaration →:

data type var-name [Expression]; a Particular element of the array.

data type → is type of element to be stored in array

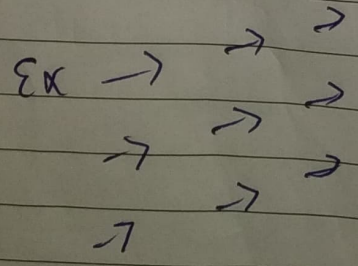
```
int a[5];
```

var-name → simple variable name



The Expression / Subscript specifies the no. of values stored in the array.

memory size of array = $5 \times 2 = 10 \text{ bytes}$



- Integer = 2 bytes
- Float = 4 bytes
- char = 1 bytes

memory size of array / space allocated to array = size of array * size of array

Index of first element $\rightarrow 0, 1, 2$

Index of last element 18

\rightarrow size of ~~array~~ element $- 1$

042

Ex \rightarrow int arr[2]

arr[0] = 5

arr[1] = 10

~~arr[2]~~

\rightarrow Index of last element

$$\rightarrow 2 - 1 = 1$$

Array num size =

Size = (upperbound - lowerbound)

+ 1

\rightarrow upperbound
arr[0]

arr[1] \rightarrow upperbound

\Rightarrow

$$\text{Size} = (1 - 0) + 1$$

$$= 1 + 1 = \underline{\underline{2}}$$

Advantages \rightarrow :

\rightarrow we can directly access the data values from array using index number.

\rightarrow we can use an array to graph related variable under a single name

Disadvantages \rightarrow :

(1) It's a necessity to declare in advance the amount of memory to be utilized by array.

\rightarrow Phere se hee batand Padega kii Kyaa Size hogi array ki.

Ex \rightarrow int

arr[10];

(X)

(2) It requires consecutive memory locations there may be not possible if array size is large.

(8 are formulae)
Date: \rightarrow 29/06/20

of Array
chart from JavaPoint

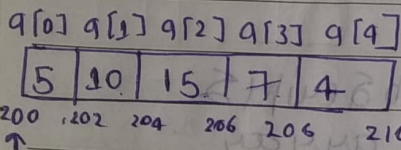
Why-linked list concept introduced?

Date: \rightarrow 01/07/20

FOR EX \rightarrow

int a[5];
Constant size

(bcz compiler allocate memory on compile time)



base address \rightarrow 5 x 2 \rightarrow 10 bytes (memory allocate)

base address

\rightarrow 5 x 2 \rightarrow 10

bytes (memory allocate)

\rightarrow Access value of index 2

Index = 2

value = a[2]

Address = Base address + Index x Size of integer

$= 200 + 2 \times 2$

$= 204$

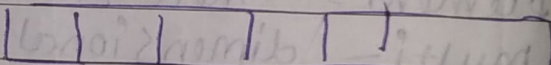
Access 3rd element of array

$= 3 - 1 = 2$

Drawback of array \rightarrow :

(i) we should know the size of array in advance.

int a[1000];



wastage of memory.

(in case float so float char)

② Contiguous memory allocation used in memory. →

Ex → 20,000 memory location available but not at one place

③ Insertion & deletion are time consumed.

↓

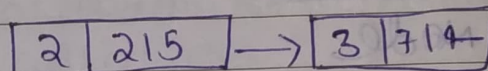
So to remove the above drawbacks of array.

Linked list concept is introduced.

Advantages →

(iii) Random memory allocation →

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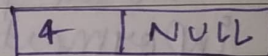


200 ↑ Head
Node

215



data → 2



714

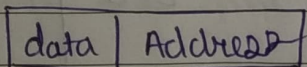
→

2 ke baad 3 aayega
Ye kaise pata chalega
ki jo 3 ki address
hai wo 2 mai
store hojayegi

Linked list →

→ We store data in the form of Node. We can't tell size ~~we~~ create one by one node to store data.

Node

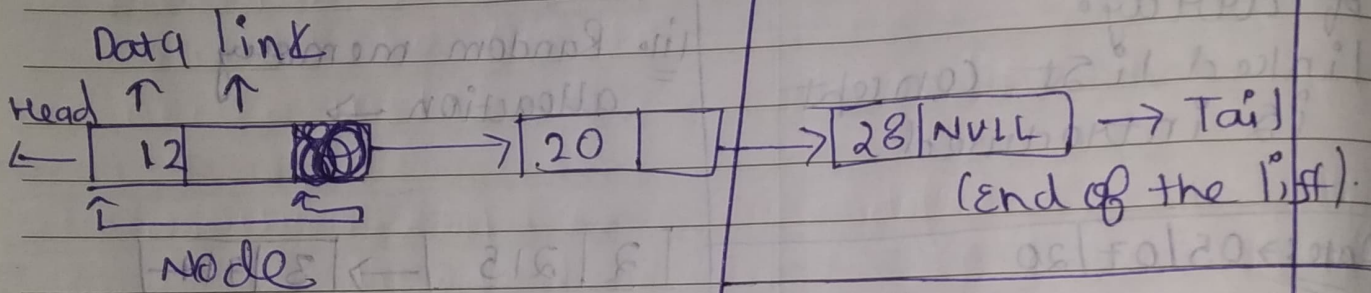


Pointer → They are used to store the address of another Node.

→ You can create any no. of nodes without giving size of list in advance.

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→ L.L. can be defined as collection of objects called nodes that are randomly stored in the memory.



→ The data field stores actual piece of information.

→ Link field is used to store ~~the address~~ the address of next node. (Pointer)

Advantages →

(i) Empty node can not be present in the linked list.

Types of L.L. →

(i) Singly L.L.

2. Doubly L.L.

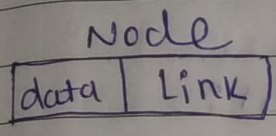
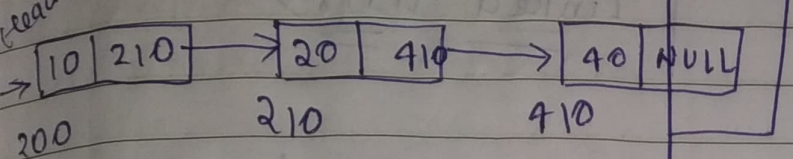
3. Circular L.L.

Singly
Circular
L.L.

Doubly
Circular
L.L.

Singly L.L. →

Ex →
read / first node



→ Collections of node which contains two fields data & Link.

→ Traversal →
One direction means you can go in forward direction. not move from 20 to 10 return. For removing this drawback L.L. introduced

Doubly L.L. →

→ It's called singly bcoz this list consists of one link to point to next node or element.

→ The first node of list is called as Head / start / first node.

→ The link field of last node is NULL which represents end of linked list.

→ Traversal is only one direction i.e. forward direction.

→ This drawback can be overcome by doubly L.L.

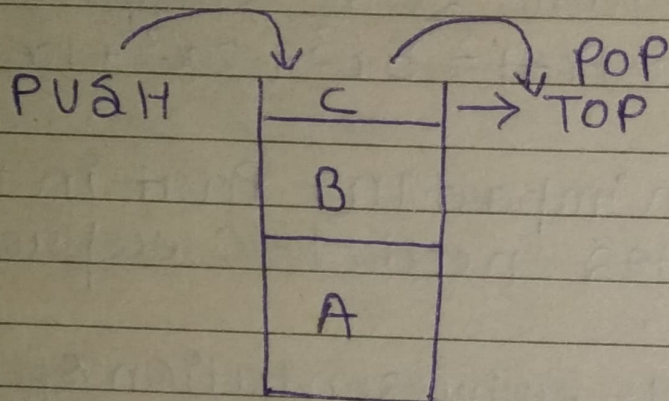
Print Reverse order \rightarrow Stack me.

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Unit \rightarrow 2

stacks



\rightarrow Non-Primitive D.O.S.
 \rightarrow linear D.O.S.

\rightarrow operation \Rightarrow Last In first out

(i). PUSH \Rightarrow Insert element

(ii). POP \Rightarrow Delete element

3. Peek or TOP \Rightarrow Return Karega last element
4° is Empty \rightarrow True \rightarrow else false.

True \rightarrow koi element Nahi hai
False \rightarrow element Hai.

Array Representation \Rightarrow :

Implementation \Rightarrow

- \rightarrow static \rightarrow array
- \rightarrow Dynamic \rightarrow Linked list

(i). linear
stack \Rightarrow is a non-primitive \forall D.O.S.

(ii). stack is a list of element in which an element may be inserted or deleted only at one ~~part~~ called top of the stack, the last ^{end} added element will be the first to be removed from the stack that is the reason the stack is also called "LIFO",

Two basic operation associate with stack \Rightarrow

(i). PUSH operation \rightarrow is used to insert an element into a stack.

(ii). POP operation \rightarrow is used to delete an element from a stack.

And ~~also~~ other operation

3^o Peek or TOP \Rightarrow Return Top element of stack.

4^o Is empty \Rightarrow Return true if stack is empty as false

Advantages ⇒

(i) Stack provides unique way to work with continuous memory.

2. Stack provide a way to access different of continuous data in a "LIFO" manner

Disadvantages ⇒

(i) Inflexible

2. Unable to copy & Paste

Applications of stacks ⇒

(i) Reverse a string

2. Recursion

3. Polish Notation

4. Expression evaluation

5. Back tracking & two operation.

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Applications of Stack \Rightarrow

(i) Polish Notation \Rightarrow

- \rightarrow Infix Notation
- \rightarrow Prefix Notation
- \rightarrow Postfix Notation

1° Infix Notation \Rightarrow Operator operand ke baad
 $x + y * m$

2° Prefix \Rightarrow Operator operand ke Bhele
 $+ xy$

3° Postfix $\Rightarrow xy +$

Remember
return when the function had to return.

Recursion make use of system stack for storing the written address of the function call.

- # Queue ⇒
- ⇒ linear D.O.S
 - Non-primitive D.O.S
 - FIFO
 - Rear (Insert → store)
 - Front (Delete → Process)

A

Queue is an linear list of elements in which deletion can take place only at one end called the "front end" and insertion can take place only at the other end called "Rear end".

Queue follows the "FIFO" methodology that is the data item stored

Static & dyn implementation of Queue

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first will be access first.

Basic operations of Queue

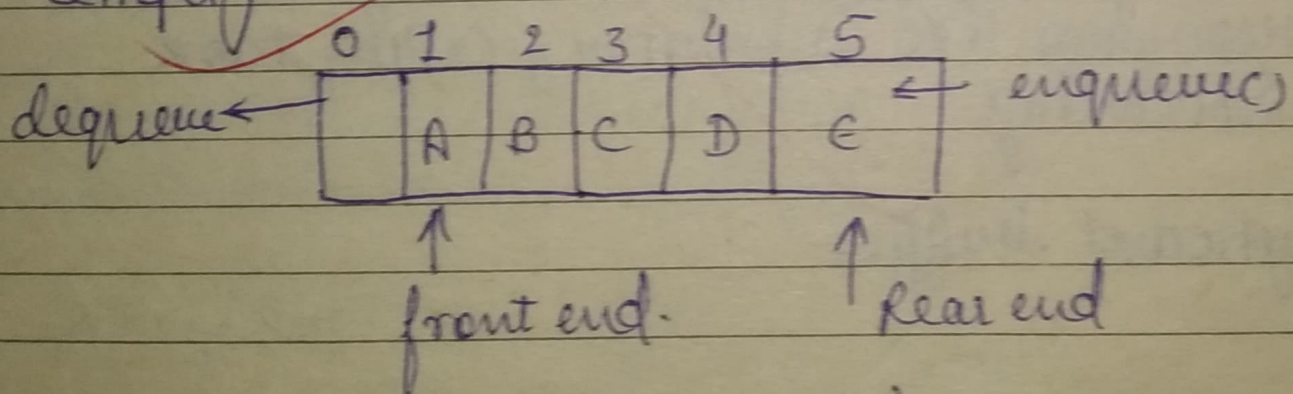
enqueue \rightarrow The process to add an element into Queue is called enqueue.

Dequeue \rightarrow The process of removal of an element from the Queue is called Dequeue.

Peek $() \rightarrow$ Gets the elements add the front of the queue without removing it.

is full $() \rightarrow$ Check if the Queue is full.

is empty $() \rightarrow$ Check if the Queue is empty.



enqueue $()$

Bin ary search
sorting
searching
Red black tree

} MST

4. Write a algo diagram for the following

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Application of Queue →

(i) Real life

(ii) Computer Science

(i) Real life ⇒

(i) Call center Phone System

(ii) Waiting in line.

2 Computer Science related ⇒:

→ Round Robin Scheduling

→ Handle of Interrupts

→ Serving a request on a single shared resource like a printer, CPU, task Scheduling

→ Keyboard buffer.

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Types of Queues

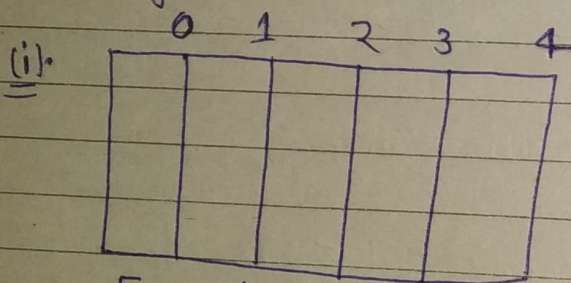
①. Simple or linear queue

②. Circular queue

③. Deque (double ended queue)

④. ~~Priority~~ Priority queue.

①. Simple or linear queue \rightarrow .



$F = -1$

$R = -1$

Queue empty

Simple queue defined the simple operation of the queue in which insertion occurs at the rear end of list & deletion occurs at the front of the list

o/

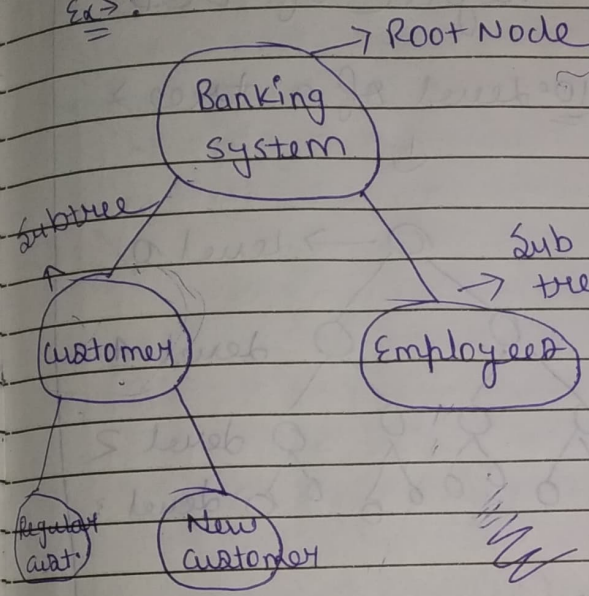
Tree

→ trees are basically used to represent the data object in hierarchical manner.

Tree Terminology: →

→ There are no. of terms associated with trees which are listed below: →

Ex: →



(i) Root Node

→ It's a unique node in the tree to which further subtrees are attached.

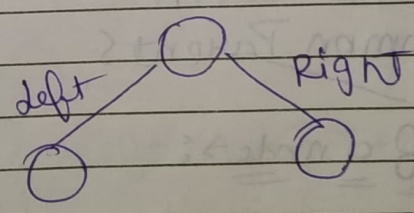
(ii) Node → Each data item in a tree is called node. It specifies the data information & links (branches) to other data items.

Definition: →

→ tree is a finite set of one or more data items (nodes) such that:

3. Parent Node → The node having further subbranches.

1. There is special data item called the root of the tree.



4. Child Node → The node having no further subtrees.

2. And it's remaining data items are partitioned into m of mutually exclusive subsets, each of which is itself a tree. They are called subtree.